

Compositional resource theories of coherence

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Quantum coherence is one of the most important resources in quantum information theory. Indeed, preventing the loss of coherence is one of the most important technical challenges obstructing the development of large-scale quantum computers. Recently, there has been substantial progress in developing mathematical resource theories of coherence, paving the way towards its quantification and control. To date however, these resource theories have only been mathematically formalised within the realms of convex-geometry, information theory, and linear algebra. This approach is limited in scope, and makes it difficult to generalise beyond resource theories of coherence for single system quantum states. In this paper we take a complementary perspective, showing that resource theories of coherence can instead be defined purely compositionally, that is, working with the mathematics of process theories, string diagrams and category theory. This new perspective offers several advantages: i) it unifies various existing approaches to the study of coherence, for example, subsuming both speakable and unspeakable coherence; ii) it provides a general treatment of the compositional multi-system setting; iii) it generalises immediately to the case of quantum channels, measurements, instruments, and beyond rather than just states; iv) it can easily be generalised to the setting where there are multiple distinct sources of decoherence; and, iv) it directly extends to arbitrary process theories, for example, generalised probabilistic theories and Spekkens toy model—providing the ability to operationally characterise coherence rather than relying on specific mathematical features of quantum theory for its description. More importantly, by providing a new, complementary, perspective on the resource of coherence, this work opens the door to the development of novel tools which would not be accessible from the linear algebraic mind set.

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1 Introduction

The understanding and manipulation of coherence within quantum systems is central to the development of quantum technologies. More than this though, deeply understanding coherence is central to our understanding of quantum theory itself. Indeed, whilst loss of coherence in a quantum system is a critical problem in the development of quantum computers, it is essential for the emergent classical world in which we live our day to day lives [90, 91, 45]. Developing better mathematical tools for addressing quantum coherence is therefore pertinent to both the foundations of physics as well as quantum engineering.

There has been great progress recently in the field of resource theories. Both the abstract mathematical formalism [43, 11, 23, 30, 17] as well as important concrete examples have been developed and applied to a broad range of problems [88, 71, 84, 63, 17, 81, 29, 68, 61]. The key benefit of the resource-theoretic perspective on quantum theory is that it takes qualitative ideas such as entanglement [44, 66], asymmetry [2], coherence [84], nonlocality [88] and so on [27], and turns them into quantitative studies. With such quantification comes the ability to optimise use of these resources and so efficiently manipulate them for the sake of performing practical tasks. In its infancy the study of resource theories was rather piecemeal. Recently however, a more systematic study has begun and unified frameworks [23, 30] have been developed for studying the abstract structure of resource theories. This has the added benefit that ideas developed in one resource theory can in a straightforward manner be ported over to another—we are no longer continually reinventing the wheel.

One of the key developments on the abstract formalism of resource theories has been the work of [23]. This utilises the mathematics of category theory and process theories [25, 73] to study resource theories in quantum theory and beyond. A key benefit of this approach is that it is intrinsically and explicitly compositional in nature. In particular, the mathematics of category theory comes with a natural diagrammatic representation which makes reasoning intuitive and greatly simplifies many complex calculations. For instance, this diagrammatic reasoning has been used to provide a natural environment for describing measurement-based quantum computation [20], quantum teleportation [19], and quantum secret sharing [21] among many others [25].

The study of resource theories of coherence has, to date, largely taken place from a more convex-geometric perspective, utilising tools such as majorization [62, 64] and Schur-convex functions [62, 35]. For example see the works of [87, 1, 10, 57, 89, 16]. However, this approach is somewhat limited in scope, and makes it difficult to generalise beyond resource theories of coherence for single system quantum states. Nowhere is understanding the manner in which multisystem, channel, and measurement coherence manifests as a resource more important than in a large-scale quantum computer, where multipartite systems evolve through the application of channels to eventually be measured. Hence understanding how coherence in such settings can be controlled and quantified—that is, used as a resource—is vital for achieving the promise of large-scale quantum computing.

In this work we take a complementary perspective and work entirely process-theoretically, equivalently we work entirely categorically, that is to say, diagrammatically. This has the benefit that it unifies many existing approaches. More importantly, it provides a general treatment of resources theories of coherence in a compositional multi-system or multipartite setting, and also generalises immediately to the study of coherence of measurements, channels, general quantum instruments and beyond. Moreover, the formalism can also be developed to allow for multiple distinct sources of decoherence for the same system. Finally, it naturally extends to the study of coherence within alternative physical theories

[32, 33], such as generalised probabilistic theories [38, 8, 39, 13] and epistemically restricted theories [82, 22, 3, 83]. This last point is useful as it provides the ability to characterise coherence in operational and physical terms, rather than via specific features of the mathematical formalism of quantum theory—such as Hilbert spaces and complex numbers. There is some existing work on studying resource theories in generalised theories such as [46, 14, 15, 86, 4] and this general approach has also been of use to deepen our understanding of computation [52, 47, 9, 31, 6, 54, 55, 50, 51], cryptography [77, 74, 48, 78, 7, 5, 53, 49] and much more.

More importantly, however, than the application to specific alternative theories, is that this general approach demonstrates that the basic essence of coherence can be captured in terms of composition of processes, and hence, that the information-theoretic approach should not be taken to be primitive.

The remainder of this paper will be structured as follows. In section 3 we introduce a process-theoretic way of understanding decoherence. In section 4 we introduce various resource theories of coherence which we can formulate from this compositional perspective. In this section we in particular: develop a resource theory for multi-partite scenarios; provide a unified treatment of speakable and unspeakable coherence; and generalise to the case in which there are multiple possible sources of decoherence. Finally we discuss these results in section 5 and point out a wide range of topics for future work. Before this however we will provide some basic background material on process theories (sec. 1.1) and resource theories (sec. 2 and app. A) which will be essential for understanding the rest of the paper.

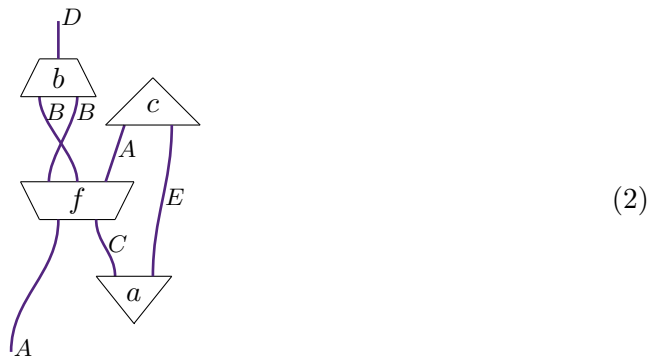
1.1 Process theories

We will provide a brief introduction to process theories here, for a more detailed introduction see, for example, [75, Section 2] or [24, 25, 73].

A process theory, \mathcal{P} , is defined by a collection of *processes*, e.g.:

$$\begin{array}{c} |B \quad |B \quad |A \\ \hline f \\ \hline |A \quad |C \end{array} \in \mathcal{P} \quad (1)$$

is a process with *input systems* A and C and *output systems* B , B and A ¹. We interpret these systems as representing physical systems or degrees of freedom, the processes then represent how these physical systems evolve or are transformed. The defining feature of a process theory is that such processes can be wired together to make *diagrams*, e.g.:



¹The labels, A , B , C , ..., on the systems are often called the system *type*.

Moreover, we can also have both classical inputs and classical outputs, in which case the input controls which (non-destructive) measurement is occurring, a.k.a. an instrument, which we denote as:


(8)

Example 1 (Quantum theory). The key example of interest for us here is quantum theory, see [25] for a detailed presentation of quantum theory as a process theory. In short however, systems in quantum theory correspond to Hilbert spaces $\mathcal{H}, \mathcal{K}, \dots$ which, for simplicity, we take to be finite dimensional. Processes with input \mathcal{H} and output \mathcal{K} correspond to completely positive trace non-increasing (CPTNI) maps from $\mathcal{B}[\mathcal{H}]$ to $\mathcal{B}[\mathcal{K}]$. If there is no input (or output) system then we take the processes to correspond to be CPTNI maps from (or to) $\mathcal{B}[\mathbb{C}]$, in particular this means that processes with no input correspond to (subnormalised) density matrices, those with no outputs to POVM elements and those with neither to probabilities. Composite systems are represented by the tensor product of the Hilbert spaces $\mathcal{H}\mathcal{K} = \mathcal{H} \otimes \mathcal{K}$. The discarding map is given by the trace. Hence, in quantum theory the causal processes are those that are trace-preserving.

2 Resource theories

Here we provide a brief introduction to the process-theoretic approach to resource theories, as originally introduced in Ref. [23]. A more detailed introduction to this framework is presented in App. A.

A resource theory specifies a set of resources and how they can be composed and freely converted. This idea, once appropriately formalised, is actually mathematically equivalent to the notion of a process theory that we have just introduced, however, rather than interpreting systems as being physical systems and processes as their evolution, we instead interpret systems as being resources, and processes as ways in which we can freely transform them. A free resource is then one that can be constructed by some free transformation ‘out of nothing’.

In Ref. [23] the authors show that there is, in some sense, a more fundamental structure which underpins a resource theory, which is known as a *partitioned process theory*. Specifically, given a process theory \mathcal{P} , in order to construct resource theories we must identify a sub-process theory of free processes $\mathcal{P}^{\text{free}}$ – the idea being that these are the physical processes which we can freely implement with no cost. For example, in constructing the resource theory of entanglement we may identify the LOCC operations as being free. This subset of free processes should be closed under composition, that is, if we can freely implement f and freely implement g then we should be able to freely do a composite of f and g . Hence, rather than being an arbitrary subset of the processes, this subset must be a sub process theory. Such a pair $(\mathcal{P}, \mathcal{P}^{\text{free}})$ is known as a *partitioned process theory*.

Given a partitioned process theory $(\mathcal{P}, \mathcal{P}^{\text{free}})$ Ref. [23] shows that there are various different resource theories, \mathcal{R} , which can be constructed. Roughly speaking, this consists of identifying the sorts of processes within \mathcal{P} which you want to consider to be resources. For example, do we want to think of just the states in \mathcal{P} as being resources, or do we want to think of the channels as being resources? There are various constructions of this sort which are discussed in Ref. [23], and we will formally introduce in the App. A. Note, however, that the precise construction that should be used to define the resource theory

will depend on exactly what one chooses to designate as a resource. In this sense, the partitioned process theory is really the more fundamental structure, and the particular resource theory that one constructs is situation dependent.

In the study of resource theories we are often not interested in the details of how we transform one resource into another (as is described by \mathcal{R}), but simply in whether or not there exists some way to freely transform one resource into another. That is, we often work simply with the *resource preorder*, R , denoted:

$$s \succ r \tag{9}$$

for certain pairs of resources $s, r \in R$. This preorder \succ means that s can be freely converted into r , and hence, that s is a more valuable resource than r . This resource preorder can be immediately obtained from a given resource theory simply by ignoring some of its structure.

It is therefore important to see that, once the partitioned process theory has been defined, the rest is completely mechanical: once we know what resources we're interested in then we can construct the relevant resource theory and so we have (in principle at least) a characterisation of this preorder. The challenge for defining interesting resource theories in practice, therefore, boils down to defining interesting partitioned process theories. Hence, this is going to be our main focus in this paper.

3 Decoherence

To understand coherence process-theoretically we will first introduce a process-theoretic understanding of decoherence. This will be crucial for our development of resource theories of coherence in section 4. This formalism has been developed for both quantum and more general theories in the following papers [26, 56, 67, 72, 33, 42, 76, 70, 41], we will again not go into the full details of this but instead just present the fragment necessary for this paper.

Decoherence is generally an extremely complicated physical process involving an interaction of our system of interest A with some uncontrolled environmental system, E , which we then lose access to. However, whilst the underlying mechanism is complex, the effective evolution of the system of interest is much simpler. Indeed, the essence of decoherence can be captured by the following simple definition.

Definition 1 (Decoherence process). A decoherence process for a system A , denoted:

$$\begin{array}{c} |^A \\ \diamond \\ |_A \end{array} \tag{10}$$

is a causal process, i.e.:

$$\begin{array}{c} \text{---} \\ |^A \\ \diamond \\ |_A \end{array} = \begin{array}{c} \text{---} \\ |_A \end{array} \tag{11}$$

which is idempotent, i.e.:

$$\begin{array}{c} |^A \\ \diamond \\ \diamond \\ |_A \end{array} = \begin{array}{c} |^A \\ \diamond \\ |_A \end{array} \tag{12}$$

This definition is motivated by the fact that decoherence should be causal, in that it is something that happens deterministically to our system rather than as a particular outcome of some measurement. Secondly, decoherence should be idempotent, as once coherence has been lost decohering again should have no further impact on the system. In general there are many different decoherence processes for a given system. Suitably choosing the decoherence processes is therefore an important part of defining a resource theory of coherence. An alternative perspective, as suggested by one of the referees, is that such processes can be viewed as the physically realisable resource destroying maps [59].

Note that there are always trivial decoherence processes for any process theory, namely the identity process:

$$\left| A \right. \quad (13)$$

these are both idempotent and causal so constitute (trivial) decoherence processes. Clearly then, the interesting resource theories are to be found by considering non-trivial decoherence maps.

Example 2 (Decoherence processes in quantum theory). Let us briefly return to quantum theory to illustrate this with a key example. Consider a qubit, then an example of a decoherence process is the completely dephasing map, \mathcal{D} :

$$\mathcal{D}(\rho_A) = \sum_{i=0}^1 |i_A\rangle\langle i_A| \rho_A |i_A\rangle\langle i_A| \quad (14)$$

however, as discussed in Ref. [63], even just for a qubit there are qualitatively different forms that decoherence could take. For example, strictly, by our definitions, the identity map is a decoherence process, albeit an extremely uninteresting one. Moreover, any discard-prepare channel is a valid decoherence process, again however one that will not lead to particularly interesting resource theories. In the end what can be shown [26] is that the qualitatively different forms of decoherence can be classified as the sub C^* -algebras of the quantum system. Intuitively, these C^* -algebras simply correspond to the different block-diagonal forms that can be embedded within the quantum system, for example, if we consider a qutrit then one C^* -algebra is the set of 3×3 density matrices, one is given by density matrices 2×2 block and a 1×1 block, and the final one is given by the completely diagonal density matrices. Our notion of decoherence here is therefore substantially more general, even for quantum theory, than the traditionally studied case of the completely dephasing map (see, e.g., [84]).

These decoherence processes are actually all that we need to define certain resource theories of coherence—in particular those representing “unspeakable coherence” as coined in [63]. However, for other resource theories of coherence we will need a little more information about the decoherence mechanisms that we are interested in. Specifically, this will be necessary for representing “speakable coherence” [63].

In particular, we need to specify not just the decoherence process, but a closely related process that we will term the *decoherence mechanism*². Intuitively, a decoherence mechanism specifies the resultant state of the environmental system in addition to the evolution of the system of interest—this is important to study how information about the system of interest leaks into the environment.

²Previously known as preleaks in [72] but this terminology is not so useful here.

Definition 2 (Decoherence mechanism). A decoherence mechanism is a process denoted

$$\begin{array}{c}
 |A\rangle \\
 \diamond \\
 |A\rangle
 \end{array}
 \begin{array}{c}
 |E\rangle \\
 \curvearrowright
 \end{array}
 \quad (15)$$

which induces a decoherence process on the system A when the environment system E is discarded, i.e.:

$$\begin{array}{c}
 |A\rangle \\
 \diamond \\
 |A\rangle
 \end{array}
 \begin{array}{c}
 |E\rangle \\
 \curvearrowright
 \end{array}
 =
 \begin{array}{c}
 |A\rangle \\
 \diamond \\
 |A\rangle
 \end{array}
 \quad (16)$$

Note that given a particular decoherence process there will, in general, be many decoherence mechanisms which could induce it. For example, the decoherence process itself can actually be viewed as a decoherence mechanism, just one where the system E is the trivial ‘no system’ system. Part of the challenge of constructing interesting resource theories of coherence is therefore to pick suitably interesting decoherence mechanisms.

Example 3 (Decoherence mechanisms in quantum theory). To give one example, a decoherence mechanism, \mathcal{B} , that could induce the completely dephasing map, \mathcal{D} , could be the following map from a qubit A to itself together with an environmental qubit E :

$$\mathcal{B}(\rho_A) = \sum_{i=0}^1 |i_A i_E\rangle \langle i_A| \rho_A |i_A\rangle \langle i_A i_E| \quad (17)$$

Typically decoherence is viewed as a process that happens to states, that is, given some state s of system A , if we have some decoherence process for A then we can define the decohered state as:

$$\begin{array}{c}
 |A\rangle \\
 \diamond \\
 |A\rangle \\
 \triangle \\
 s
 \end{array}
 \quad (18)$$

and hence we can define the *decohered state space* as the collection of all such states. However, unlike in many other approaches to studying decoherence we can also define decoherence of general processes, that is, given some process $f : A \rightarrow B$, and a decoherence mechanism for both A and B then we can define the decohered process as those that are pre and post-composed with decoherence processes:

$$\begin{array}{c}
 |B\rangle \\
 \diamond \\
 |B\rangle \\
 \square \\
 f \\
 |A\rangle \\
 \diamond \\
 |A\rangle
 \end{array}
 \quad (19)$$

and so we can define the *decohered process space* as the collection of all such processes. Due to idempotence of the decoherence processes, we can equivalently define the decohered processes as being those that satisfy:

$$\begin{array}{c}
 |B\rangle \\
 \square \\
 f \\
 |A\rangle
 \end{array}
 =
 \begin{array}{c}
 |B\rangle \\
 \diamond \\
 |B\rangle \\
 \square \\
 f \\
 |A\rangle \\
 \diamond \\
 |A\rangle
 \end{array}
 \quad (20)$$

or, again equivalently, the pair of constraints that:

$$\begin{array}{c} |B \\ \boxed{f} \\ |A \\ \diamond \\ |A \end{array} = \begin{array}{c} |B \\ \boxed{f} \\ |A \end{array} = \begin{array}{c} |B \\ \diamond \\ |B \\ \boxed{f} \\ |A \end{array} \quad (21)$$

Going further however, we can define decoherence for a process theory \mathcal{P} itself. To do so, for each system A , we choose a decoherence process, we denote this collection as:

$$\text{dec} := \left\{ \begin{array}{c} |A \\ \diamond \\ |A \end{array} \mid A \in \mathcal{P} \right\} \quad (22)$$

Note that if we have classical systems in our process theory then we will usually want to interpret these as having trivial decoherence, that is, where the chosen decoherence process is simply the identity.

We can then define the set of decohered processes in \mathcal{P} , which we denote as \mathcal{P}^{dec} , as:

$$\mathcal{P}^{\text{dec}} := \left\{ \begin{array}{c} |B \\ \boxed{f} \\ |A \end{array} \in \mathcal{P} \mid \begin{array}{c} |B \\ \boxed{f} \\ |A \end{array} = \begin{array}{c} |B \\ \diamond \\ |B \\ \boxed{f} \\ |A \end{array} \right\} \quad (23)$$

This is the minimal structure that we require to construct resource theories. However, it will commonly be the case that we want \mathcal{P}^{dec} to itself be a process theory, that is, it should be closed under wiring together processes. Intuitively, if we compose two processes which don't possess any coherence, then the resulting process also shouldn't have any coherence. That is, the act of composition doesn't spontaneously create coherence. Additionally, such a feature would be desirable, for instance, if we want to view decoherence as describing how an entire theory emerges from another, for example, to describe how classical theory emerges from quantum theory via decoherence.

To ensure that \mathcal{P}^{dec} is closed under composition we must impose certain constraints on the choice of the decoherence processes. The specific constraint to impose will depend on the specific situation that one has in mind, however, there are two choices that seem to crop up naturally.

Example 4. The easiest way to enforce such a compositionality constraint is to simply demand, for all systems A and B , that:

$$\begin{array}{c} |A \\ \diamond \\ |A \end{array} \begin{array}{c} |B \\ \diamond \\ |B \end{array} = \begin{array}{c} |A \\ \text{---} \\ |A \end{array} \begin{array}{c} |B \\ \text{---} \\ |B \end{array} \quad (24)$$

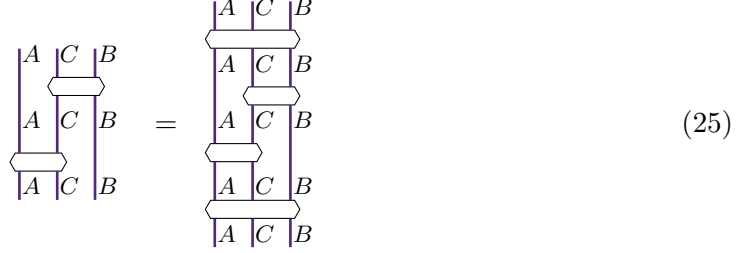
This would be a suitable choice, for example, if we view these decoherences as occurring in a quantum computer where it is typically assumed that noise is uncorrelated.

Lemma 1. If the decoherence process for composite system is the composite of the decoherence processes for the components, then \mathcal{P}^{dec} is closed under composition.

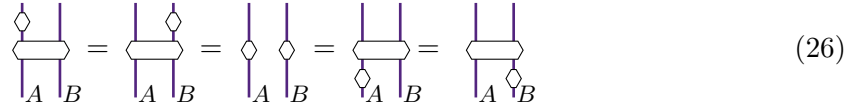
Proof. See appendix B □

This constraint on decoherence for composite systems may often be the appropriate choice³, however, there are cases where this is too restrictive. For example, if we want to consider decoherence arising from loss of some global reference frame then we shouldn't lose the correlations between local systems. In this case then there is a set of weaker conditions which are satisfied which still ensure that \mathcal{P}^{dec} is closed under composition.

Example 5. Another way to enforce compositionality of \mathcal{P}^{dec} is by demanding i) that the composition of decoherence processes are decohered processes, i.e. that for all systems A, B and C we have:



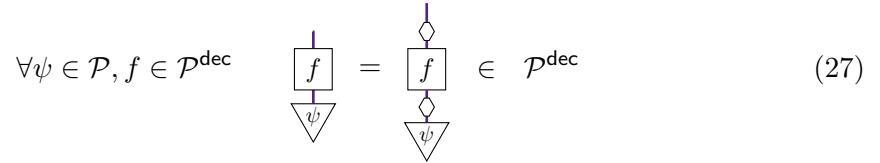
and, ii) that if we compose a global with a local decoherence we get a pair of local decoherences, i.e. that for all systems A and B we have:



Lemma 2. The conjunction of these two constraints (equations 25 and 26) imply that \mathcal{P}^{dec} is closed under composition.

Proof. See appendix C □

If one were to pick the decoherence processes such that \mathcal{P}^{dec} is indeed closed under composition, then this defines a partitioned process theory⁴. One might then be tempted to use this as the definition of a partitioned process theory from which to construct resource theories, that is, defining a partitioned process theory by $(\mathcal{P}, \mathcal{P}^{\text{free}} := \mathcal{P}^{\text{dec}})$ ⁵. However, this is far too restrictive a definition of free processes. Specifically, we can see that the resource theory of states has an extremely boring preorder. To see this, note that if take the decohered processes as being the free processes then any free process, f , maps any arbitrary (potentially non-free) state ψ in \mathcal{P} into the free (i.e. decohered) partition:



Hence, the preordering of the resource theory of states for this partitioning is too coarse grained to capture interesting phenomena. We will therefore, in the following section, explore how to construct interesting resource theories of coherence.

³Indeed, it is by far the most commonly studied situation in the literature.

⁴Strictly this is not a partitioned process theory in the sense of [23] as the free set does not include identities. There are various ways around this technicality, for instance, we could simply add identity processes to the free set.

⁵Indeed such a definition has been studied in [69].

4 Resource theories of coherence

The above mathematical concepts that we have introduced provide us with a tool-box for constructing resource theories of coherence. The specific resource theory that one constructs from these will really depend on the specific application that one has in mind. In this section we will therefore illustrate this toolbox with some example resource theories, however, we stress that this is not intended to be exhaustive and should be viewed more as a pedagogical introduction to these tools.

In doing so, we encounter some of the resource theories previously discussed in the literature—such as Decoherence Invariant Operations (DIO), and Translationally Invariant Operations (TIO) [84]—as special cases. Moreover, we explicitly construct a resource theory of coherence for multipartite systems, the free processes of which we term Complete Decoherence Invariant Operations. Importantly, we show that the constraints on free processes in this theory are not implied by any previous theory in the literature.

There are (at least) two qualitative choices to be made when deciding which partitioned process theory to construct, firstly, there is the question of whether we want the coherence to be speakable or unspeakable—as defined in [63]. Intuitively, consider the two quantum states $\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$ and $\frac{1}{\sqrt{2}}(|0\rangle + |2\rangle)$ of a qutrit. If these are treated as equivalent resources then we are dealing with *speakable* coherence, where all that matters is whether there is a superposition, not which subspaces are involved in the superposition (such as would be the case in quantum computation). On the other hand, if these are inequivalent resources then we are dealing with *unspeakable* coherence, where which subspaces appear in superposition is relevant (such as would be the case if we were using these states for metrology). Secondly, there is the question of how to deal with composite systems, namely, is coherence a local or a global phenomena⁶. In both cases we can define “intermediate” cases as well but the two extremes are the most interesting. Then, when moving from the partitioned process theory to a resource theory there is the question of what is the resource, are they states, general processes, collections of processes or something more complicated such as quantum combs?

Regardless of all of these choices, we will demand that our resource theories must abide by one basic principle:

Principle 1. *Free processes must preserve the set of decohered processes.*

Intuitively, free processes cannot create coherence spontaneously, and so should map any process lacking coherence to another. We will show how to formalise this intuition in the following subsection. Interestingly, we will then see that some free processes defined in the literature, such as Maximally Incoherent Operations (MIO) [84], fail to satisfy this principle.

4.1 Single system speakable coherence

We will start with a very simple example which will help to illustrate the basic concepts: a process theory with only a single system (and the trivial system) and, hence, there will be just one relevant decoherence map. From this, we will construct the resource theory of states. This is not a particularly physically well motivated process theory as any physical theory will have non-trivial composite systems, however, we can still see many of the important features appearing in this simplified setting.

⁶This depends on what we view as being the source of decoherence, in general within quantum computation we assume that errors are local, and so systems decohere locally. On the other hand, if we view decoherence as arising from the loss of some global reference frame then it makes sense to define decoherence as a global phenomena.

Given a process theory \mathcal{P} with a single system A we can pick a particular decoherence process:

$$\begin{array}{c} A \\ | \\ \diamond \\ | \\ A \end{array} \quad (28)$$

we can then, as before, define the decohered subtheory as the collection of processes:

$$\mathcal{P}^{\text{dec}} := \left\{ \begin{array}{c} A \\ | \\ \diamond \\ | \\ A \\ \triangleleft \\ s \end{array}, \begin{array}{c} A \\ | \\ \diamond \\ | \\ A \\ \square \\ f \\ | \\ A \\ \diamond \\ | \\ A \end{array}, \begin{array}{c} \triangleleft \\ e \\ | \\ \diamond \\ | \\ A \end{array} \right\} \quad (29)$$

Now, a minimal requirement for the free processes must be that they preserve the set of decohered processes. For example, a free process $g : A \rightarrow A$ must satisfy:

$$\begin{array}{c} | \\ \square \\ | \\ \diamond \\ | \\ \triangleleft \\ s \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \square \\ | \\ \diamond \\ | \\ \triangleleft \\ s \end{array} \quad (30)$$

that is, it must preserve the set of free states. This example is commonly taken as the only constraint on free processes, corresponding, at least in the quantum case, to the free processes known as MIO in the literature [84]. However, now that we are taking a process-theoretic viewpoint we can see that this is not a good choice for free operations, as this is just one way that g could be combined with a decohered process. Another way that it can be combined with a decohered process is to compose it with a decohered effect, in which case we additionally obtain the constraint:

$$\begin{array}{c} \triangleleft \\ e \\ | \\ \diamond \\ | \\ \square \\ | \\ \diamond \\ | \\ \square \\ | \\ \diamond \\ | \\ \triangleleft \\ s \end{array} = \begin{array}{c} \triangleleft \\ e \\ | \\ \diamond \\ | \\ \square \\ | \\ \diamond \\ | \\ \square \\ | \\ \diamond \\ | \\ \triangleleft \\ s \end{array} \quad (31)$$

Indeed, if we consider all such possibilities we can show that our free processes must satisfy:

Lemma 3. That a process $g \in \mathcal{P}^{\text{free}}$ preserves the decohered subtheory \mathcal{P}^{dec} is equivalent to the condition:

$$\begin{array}{c} A \\ | \\ \diamond \\ | \\ A \\ \square \\ g \\ | \\ A \end{array} = \begin{array}{c} A \\ | \\ \square \\ | \\ A \\ \diamond \\ | \\ A \end{array} \quad (32)$$

Proof. To see this it suffices to consider composition of a free process with the decohered process which is just the decoherence process itself. It is simple to see that the decoherence process is indeed in \mathcal{P}^{dec} due to idempotence of the decoherence process. Pre-composing g with the decoherence process must give another process h in the decohered theory,

$$\begin{array}{c} | \\ \square \\ | \\ \diamond \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \square \\ | \\ \diamond \end{array} \quad (33)$$

then, using idempotence of decoherence, it is simple to show that this gives us

$$\begin{array}{c} \text{---} \\ | \\ \square \text{ } g \\ | \\ \diamond \\ | \\ \text{---} \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \square \text{ } h \\ | \\ \diamond \\ | \\ \text{---} \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \diamond \\ | \\ \square \text{ } h \\ | \\ \diamond \\ | \\ \text{---} \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \text{---} \\ | \\ \square \text{ } g \\ | \\ \diamond \\ | \\ \text{---} \end{array} \quad (34)$$

Similarly, we can ask that the post-composing g with the decoherence map must give another process in the decohered theory which allows us to show that:

$$\begin{array}{c} | \\ \diamond \\ | \\ \square \text{ } g \\ | \\ \text{---} \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \square \text{ } h' \\ | \\ \diamond \\ | \\ \text{---} \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \diamond \\ | \\ \square \text{ } h' \\ | \\ \diamond \\ | \\ \text{---} \end{array} = \begin{array}{c} | \\ \diamond \\ | \\ \text{---} \\ | \\ \square \text{ } g \\ | \\ \diamond \\ | \\ \text{---} \end{array} \quad (35)$$

the conjunction of these two gives the desired constraint. What remains is to show that this implies that any decohered process is preserved by g . This follows immediately from the above constraint together with eq. (23) in the definition of the decohered subtheory. \square

Such a constraint on the free processes is known as DIO in the literature [84]. That is, we have shown that

$$\mathcal{P}^{\text{free}} \subseteq \text{DIO} \quad (36)$$

Therefore, any proposal for a resource theory of coherence which use a set of free processes larger than DIO are inconsistent with our basic principle of free processes (Principle 1), namely, they map decohered processes to processes which are not decohered⁷.

If we work with the definition that $\mathcal{P}^{\text{free}} := \text{DIO}$ then we can, using the methods discussed in the earlier sections, define from this a resource theory of coherence of states (or measurements, or channels), and from that obtain a preorder amongst the states (or measurements, or channels) which fully quantifies their coherence.

One should note at this point that one of the major differences of this approach to more conventional approaches to the study of coherence is the basic building blocks that we use. Typically, the starting point of the study of coherence is to pick a preferred basis, and from there to define coherence relative to that basis. This then demands that we explain why we chose that particular basis, what made it distinct from any other basis? Here we do not pick a basis at all. Instead, we pick a particular decoherence map. This decoherence map is given to us by the physics of the situation which tells us how our system decoheres. We could, if we wanted, characterise the decoherence via process tomography, but this is not necessary to prove general results that would hold regardless of the precise nature of decoherence. Such a decoherence process may in turn pick out a preferred basis (for example, if it is a totally dephasing map in quantum theory), however, importantly this preferred basis is picked out by the physical decoherence process rather than being something that we had to choose at the start.

Let us now illustrate these definitions with a simple example.

Example 6 (Resource theory of coherence of qubit states). Consider a process theory describing a single qubit, with the decoherence map given by $\mathcal{D}(\cdot) = \sum_{i=0}^1 |i\rangle\langle i| \cdot |i\rangle\langle i|$. The decohered processes are then: states are of the form $\sum_i p_i |i\rangle\langle i|$ where $p_i \in [0, 1]$ and

⁷This is closely related to a result in [34] in the context of resource destroying maps in quantum theory which similarly singles out DIO as the largest relevant set of free processes.

$\sum_i p_i = 1$; effects are of the form $\sum_i r_i |i\rangle\langle i|$ where $r_i \in [0, 1]$; and, general transformations are of the form $\sum_{ij} s_{ij} |i\rangle\langle j|$ where s_{ij} is a stochastic matrix.

The free processes are then those that preserve the set of decohered processes, or equivalently, that commute with the dephasing map \mathcal{D} . That is: for states and effects these are simply the decohered states and effects; whilst for transformations they are those such that satisfy $\mathcal{D} \circ \mathcal{E} = \mathcal{E} \circ \mathcal{D}$.

We can then explicitly construct the relevant resource theory for states, see Ex. 7 in App. A. Specifically, the resources correspond to arbitrary qubit states ρ , and there is a transformation from resource ρ to resource σ if and only if there is a free process mapping ρ to σ :

$$\exists \boxed{\mathcal{E}} \begin{array}{c} \sigma \\ | \\ \rho \end{array} \iff \exists \mathcal{E} \text{ s.t. } \mathcal{D} \circ \mathcal{E} = \mathcal{E} \circ \mathcal{D} \text{ and } \boxed{\mathcal{E}} = \begin{array}{c} | \\ \sigma \\ \hline \rho \end{array} \quad (37)$$

If within the resource theory there exists such a process $\mathcal{E} : \rho \rightarrow \sigma$ then we write that $\rho \succ \sigma$, or in words, that ρ is at least as resourceful as σ .

The fact that free processes must be DIO does not, however, imply that taking a smaller set of free processes cannot be justified. There are two instances in which this may be well motivated. The first, which we will return to later, is when we investigate composite systems, the second, which we consider next, is when we deal with unspeakable coherence.

4.2 Single system unspeakable coherence

We can deal with unspeakable coherence within our toy example of a process theory with a single system. To do so, rather than just picking a particular decoherence process for the system, we must also pick a suitable decoherence mechanism, that is:

$$\begin{array}{c} A \\ | \\ \diamond \\ | \\ A \end{array} \quad \text{and} \quad \begin{array}{c} A \\ | \\ \diamond \\ | \\ A \end{array} \begin{array}{c} E \\ | \\ \hline \end{array} \quad (38)$$

such that

$$\begin{array}{c} A \\ | \\ \diamond \\ | \\ A \end{array} = \begin{array}{c} A \\ | \\ \diamond \\ | \\ A \end{array} \begin{array}{c} E \\ | \\ \hline \end{array} \quad (39)$$

the natural generalisation of the definition of the DIO processes is then given by the following condition:

$$\begin{array}{c} A \\ | \\ \diamond \\ | \\ A \end{array} \begin{array}{c} E \\ | \\ \hline \end{array} = \begin{array}{c} A \\ | \\ \diamond \\ | \\ A \end{array} \begin{array}{c} E \\ | \\ \hline \end{array} \quad (40)$$

This condition states that not only do we require that free processes commute with decoherence, but also that the effect of decoherence on the environment is left invariant by the free processes. This immediately implies the condition of DIO as can be seen by discarding

system E in the above equation:

$$(41)$$

Therefore this set of processes is contained within DIO. Precisely how restrictive this condition is will depend on the precise nature of this decoherence mechanism. It could be a trivial restriction, for example, in the case that

$$(42)$$

for some normalised state s , this clearly reduces to the definition of DIO. However, in other cases this is much more restrictive. For example, in quantum theory, if we take the decoherence mechanism to be

$$\mathcal{B}(\rho_A) = \sum_{i=0}^1 |i_A i_E\rangle \langle i_A| \rho_A |i_A\rangle \langle i_A i_E| \quad (43)$$

then, this corresponds to the set of processes known as TIO in the literature [84].

To see this more generally, let us sketch out how to describe decoherence as arising from the lack of access to some reference frame within quantum theory⁸. Suppose we have some group G which represents the possible orientations of the reference frame. Then we have some physical system which encodes this orientation in a set of perfectly distinguishable states:

$$(44)$$

which are perfectly distinguished by the effects:

$$(45)$$

Then, we introduce the action of the group on the physical system:

$$(46)$$

We can then write the decoherence mechanism arising from this reference frame as:

$$(47)$$

⁸This clearly applies much more broadly than just to quantum theory but would require more process-theoretic formalism than we introduce here. A formal treatment of the general case is therefore left to future work.

such that the decoherence process is given by loss of the reference frame as:

$$\begin{array}{c} \diamond \\ | \end{array} = \begin{array}{c} \overline{\overline{\overline{\diamond}}} \\ | \end{array} = \int_G d\gamma \begin{array}{c} \overline{\overline{\overline{\diamond}}} \\ | \\ R_\gamma \\ | \\ \overline{\overline{\overline{\triangle}}} \\ s_\gamma \\ | \end{array} = \int_G d\gamma \begin{array}{c} | \\ R_\gamma \\ | \end{array} \quad (48)$$

such a decoherence process is known as a G -twirling map in the literature. Recall that our constraint on free processes is that:

$$\begin{array}{c} A \quad E \\ | \quad | \\ \diamond \\ | \\ A \\ | \\ f \\ | \\ A \end{array} = \begin{array}{c} A \quad E \\ | \quad | \\ f \\ | \\ A \\ | \\ \diamond \\ | \\ A \end{array} \quad (49)$$

In the case we are considering this means that:

$$\int_G d\gamma \begin{array}{c} | \\ f \\ | \\ R_\gamma \\ | \\ \overline{\overline{\overline{\triangle}}} \\ s_\gamma \\ | \end{array} = \int_G d\gamma \begin{array}{c} | \\ R_\gamma \\ | \\ \overline{\overline{\overline{\triangle}}} \\ s_\gamma \\ | \\ f \\ | \end{array} \quad (50)$$

which implies that

$$\forall \gamma' \in G \quad \int_G d\gamma \begin{array}{c} | \\ f \\ | \\ R_\gamma \\ | \\ \overline{\overline{\overline{\triangle}}} \\ s_\gamma \\ | \\ \overline{\overline{\overline{\triangle}}} \\ e_{\gamma'} \\ | \end{array} = \int_G d\gamma \begin{array}{c} | \\ R_\gamma \\ | \\ \overline{\overline{\overline{\triangle}}} \\ s_\gamma \\ | \\ f \\ | \end{array} \quad (51)$$

and so, due to distinguishability of the reference frame states, we find that

$$\forall \gamma' \in G \quad \begin{array}{c} | \\ f \\ | \\ R_{\gamma'} \\ | \end{array} = \begin{array}{c} | \\ R_{\gamma'} \\ | \\ f \\ | \end{array} \quad (52)$$

That is, f commutes with the action of the group G – precisely the definition of TIO. Note that by varying the distinguishability of the states s_γ we can interpolate between this TIO and the DIO conditions. We no longer have a speakable vs. unspeakable dichotomy, but rather we have a spectrum between the two.

4.3 Multi-system speakable coherence

We now turn to much more interesting process theories, of which quantum theory is an example, where we have multiple different systems—we must now contend with the added difficulty (and interest!) afforded by parallel composition. This is the first time a resource theory of coherence has been developed for multiple systems, however, the multipartite case has also been studied from the perspective of combining entanglement and coherence into a single resource theory [18, 85] using restricted classes of LOCC operations..

We will continue to adhere to our basic principle (Principle 1): free processes must preserve the set of decohered processes. For example, if $h \in \mathcal{P}^{\text{dec}}$ and $g \in \mathcal{P}^{\text{free}}$, then it

must be the case that:

$$\begin{array}{c} \boxed{g} \\ \boxed{h} \end{array} \in \mathcal{P}^{\text{dec}} \quad (53)$$

However, the sorts of composites that we consider here cannot be completely arbitrary, for example, we do not expect the parallel composite of a decohered process with a free process to itself be decohered. For instance the identity process belongs to the free set but we would expect, for any $h \in \mathcal{P}^{\text{dec}}$, that:

$$\begin{array}{c} | \\ \boxed{h} \\ | \end{array} \notin \mathcal{P}^{\text{dec}} \quad (54)$$

Our key principle should therefore be formalised as follows:

Definition 3. The minimal constraint that we will impose on $\mathcal{P}^{\text{free}}$ is therefore that, if g is free then, for any $h, h' \in \mathcal{P}^{\text{dec}}$ with suitable system types we have:

$$\begin{array}{c} \boxed{g} \\ \boxed{h} \end{array} \in \mathcal{P}^{\text{dec}} \quad \text{and} \quad \begin{array}{c} \boxed{h'} \\ \boxed{g} \end{array} \in \mathcal{P}^{\text{dec}} \quad (55)$$

It is simple to show that this is equivalent to a ‘‘complete’’ notion of decoherence invariance, which we coin **cDIO**.

Definition 4. Given a set of decoherence processes **dec** we can define the set of **cDIO** processes as:

$$\text{cDIO} := \left\{ \begin{array}{c} \boxed{g} \\ |A \end{array} \begin{array}{c} |B \\ \end{array} \in \mathcal{P} \quad \left| \quad \begin{array}{c} |B \quad |C \\ \boxed{g} \\ |A \quad |C \end{array} \right. = \left. \begin{array}{c} |B \quad |C \\ \boxed{g} \\ |A \quad |C \end{array} \right\} \quad \forall C \in \mathcal{P} \quad (56)$$

where

$$\begin{array}{c} \boxed{} \\ | \end{array} \in \text{dec} \quad (57)$$

Lemma 4. A process g is in **cDIO** if and only if it satisfies the minimal constraint on free processes (def. 3). That is, def. 3 and def. 4 are equivalent.

Proof. See appendix D. □

This means that the largest set of free processes that we can consider are the **cDIO** processes, i.e. $\mathcal{P}^{\text{free}} \subseteq \text{cDIO}$.

Now, if we are in the special case, which has been most commonly studied in the literature, in which we have that global decoherence processes are just the composite of local decoherence processes, that is, as discussed in example 5:

$$\begin{array}{c} |A \quad |B \\ \diamond \quad \diamond \\ |A \quad |B \end{array} = \begin{array}{c} |A \quad |B \\ \boxed{} \\ |A \quad |B \end{array} \quad (58)$$

then $\text{cDIO} = \text{DIO}$, i.e. we can conclude that a process is cDIO simply by checking that:

$$\begin{array}{c} \diamond^A \\ | \\ \square^g \\ | \\ A \end{array} = \begin{array}{c} A \\ | \\ \square^g \\ | \\ \diamond^A \end{array} \quad (59)$$

One may therefore be interested in asking whether cDIO is a genuine restriction over DIO , or whether they are always equivalent. That is, can we find a situation in which $\text{cDIO} \subsetneq \text{DIO}$? This turns out to be the case—as we will prove via an explicit example shortly—hence, cDIO is a genuine restriction over DIO in certain circumstances.

To illustrate this, let us consider the case within quantum theory, where we are viewing decoherence as loss of a global reference system. That is, the decoherence process for a composite system is given by:

$$\begin{array}{c} | \dots | \\ | \dots | \end{array} = \int_G d\gamma \begin{array}{c} | \\ | \\ \square^{R_\gamma} \\ | \\ \square^{R_\gamma} \\ | \\ \dots \\ | \\ | \\ | \end{array} \quad (60)$$

It is simple to check that this satisfies the constraint on composition of decoherence processes described in example 5 and so that \mathcal{P}^{dec} is closed under composition. Intuitively, this form of decoherence loses any correlations to some external reference frame, but maintains correlations between subsystems. To see that this is not equivalent to DIO let us consider just a pair of qubit systems, G to be the group \mathbb{Z}_2 with representation R_γ as either the identity, $\mathbb{1}$, or a bit flip, X . Then, it is simple to compute, by considering the equation:

$$\begin{array}{c} Q_2 \\ | \\ \square^f \\ | \\ \diamond^0 \\ | \\ Q_2 \end{array} = \begin{array}{c} Q_2 \\ | \\ \diamond^1 \\ | \\ \square^f \\ | \\ \diamond^0 \\ | \\ Q_2 \end{array} \quad (61)$$

that:

$$\begin{array}{c} Q_2 \\ | \\ \square^X \\ | \\ \square^f \\ | \\ Q_2 \end{array} = \begin{array}{c} Q_2 \\ | \\ \square^f \\ | \\ \square^X \\ | \\ Q_2 \end{array} \quad (62)$$

Hence, that the cDIO transformations are those that are translationally invariant with respect to the action of \mathbb{Z}_2 . This is known to be a genuine restriction over DIO , see, for example, [63].

We have therefore found, by considering composite systems, that we are forced to revise our notions of what free processes are, even when those processes themselves do not involve composite systems. This, in retrospect is perhaps not surprising, as we know that complete positivity is a genuine restriction over positivity—the restriction from DIO to cDIO is analogous. It is therefore crucial that when constructing resource theories generally, and, in particular, for resource theories of coherence, that we do not just focus on single system phenomena but consider general forms of composition and the restrictions that they impose.

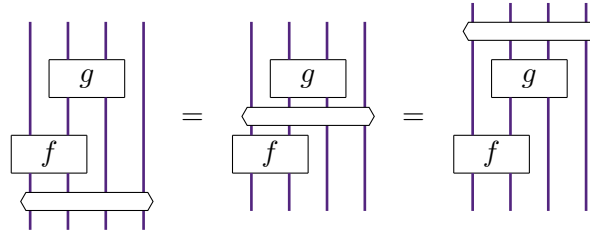
Finally, if we want to use the cDIO processes to define resource theories, that is, if we want to work with $\mathcal{P}^{\text{free}} = \text{cDIO}$ then it must be the case that this defines a partitioned process theory. In particular, we must check that cDIO processes are closed under composition.

Lemma 5. cDIO processes are closed under composition.

Proof. Consider two processes in cDIO, f and g and a general composite process:


(63)

we want to show that this too is in cDIO. The proof of this is extremely simple:


(64)

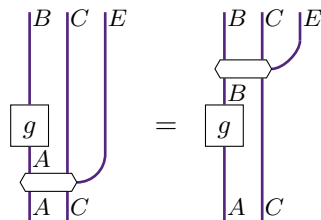
where the first equality uses the fact that f is in cDIO and the second that g is in cDIO. \square

4.4 Multi-system unspeakable coherence

Given the definition of speakable coherence above for multiple systems and of unspeakable coherence for single systems, it is simple to see how one should define speakable coherence for multiple systems. Namely, for each system A we define a decoherence mechanism:

$$\left\{ \begin{array}{c} A \\ \text{---} \\ \text{---} \\ A \end{array} \begin{array}{c} E \\ \text{---} \\ \text{---} \\ E \end{array} \middle| A \in \mathcal{P} \right\} \quad (65)$$

Then, we can define the free processes as those that commute with decoherence mechanisms in a complete sense, that is, for all systems C we demand:


(66)

Note that one limitation of this approach is that we require that each decoherence mechanism shares the same environment system E . One could avoid this at the expense of introducing encoding and decoding maps to transform between different environments. However, we can always avoid this by letting E be suitably large such that any other environment can be trivially embedded within it. Moreover, this level of generality does

not seem to add anything new conceptually so we leave its exploration to future works in which it may be relevant.

It is important to now check two things. The first check, that this set of free processes is contained within cDIO as cDIO was the minimal requirement that we wanted our free processes to satisfy. Specifically, when we define the cDIO processes relative to the decoherence processes:

$$\text{dec} := \left\{ \begin{array}{l} \begin{array}{c} |A \\ \diamond \\ |A \end{array} \\ := \begin{array}{c} |A \\ \overline{\overline{E}} \\ |A \end{array} \end{array} \middle| A \in \mathcal{P} \right\} \quad (67)$$

Secondly, we must check that processes satisfying this constraint are closed under composition such that they can be used to define a partitioned process theory. These two checks are indeed satisfied as is shown in the following simple results.

Lemma 6. The processes that commute with decoherence mechanisms must also commute with the associated decoherence processes.

Proof. The proof is trivial, simply compose the ‘commuting with decoherence mechanisms’ condition with a discarding process on the environment system:

$$\begin{array}{c} |B \quad |C \quad |E \\ \boxed{g} \\ |A \\ \diamond \\ |A \end{array} = \begin{array}{c} |B \quad |C \quad |E \\ \boxed{g} \\ |B \\ |A \end{array} \Rightarrow \begin{array}{c} |B \quad |C \\ \boxed{g} \\ |A \\ \diamond \\ |A \end{array} \overline{\overline{E}} = \begin{array}{c} |B \quad |C \quad \overline{\overline{E}} \\ \boxed{g} \\ |B \\ |A \end{array} \quad (68)$$

which immediately, given the definition of the decoherence processes, gives the desired result:

$$\begin{array}{c} |B \quad |C \\ \boxed{g} \\ |A \\ \diamond \\ |A \end{array} = \begin{array}{c} |B \quad |C \\ \boxed{g} \\ |B \\ |A \end{array} \quad (69)$$

□

Lemma 7. Processes satisfying (66) are closed under composition.

Proof. Consider two processes, f and g satisfying this constraint and a general composite process:

$$\begin{array}{c} | \\ \boxed{g} \\ | \\ \boxed{f} \\ | \end{array} \quad (70)$$

we want to show that this too satisfies this constraint. The proof of this is again extremely simple:

$$(71)$$

where the first equality uses the fact that f satisfies the constraint and the second that g does as well. \square

This therefore defines a valid partitioned process theory in which the free partitioned preserves the set of decohered processes. Hence we can, as in the other cases, use this to define resource theories of states, measurements, general processes, and so on.

Again, as we found in the single system case, we can ‘interpolate’ between speakable and unspeakable coherence by modifying the precise details of the chosen decoherence mechanisms. Intuitively, if the environment perfectly encodes the eigenspace then we will have resource theories of unspeakable coherence, and if it encodes no information then it will be speakable coherence.

Like in the single-system case, a particular instantiation of this will correspond to translationally invariant operations. However, like in the multipartite speakable case, we will find that we need to make this a ‘complete’ notion which we call **cTIO**. Specifically, given some group G and some representation on each composite system, then we define:

$$\mathbf{cTIO} := \left\{ \begin{array}{c} \begin{array}{|c} \hline B \\ \hline f \\ \hline A \\ \hline \end{array} \in \mathcal{P} \quad \left| \quad \begin{array}{|c|c} \hline B & C \\ \hline f & R_\gamma \\ \hline A & \\ \hline A & C \\ \hline \end{array} = \begin{array}{|c|c} \hline B & C \\ \hline R_\gamma & \\ \hline B & \\ \hline f & \\ \hline A & C \\ \hline \end{array} \quad \forall C \in \mathcal{P}, \gamma \in G \end{array} \right\} \quad (72)$$

Again, like in the case of **cDIO**, we find that **cTIO** can be a genuine restriction over **TIO**. However, if we are in the special case where the group representations factorise over composite systems then it will be the case that **cTIO** and **TIO** coincide. Therefore, it is imperative to specify the group representation, not just for single systems, but for composite systems as well, in order to define consistent resource theories of unspeakable coherence.

4.5 Multiple sources of decoherence

In all of the above we have considered just a single decoherence process/mechanism per system. There will be, however, situations in which we may be interested in multiple different decoherence processes/mechanisms for the same system. In this section we will briefly introduce a formalism that allows us to handle such situations. This is closely related to the categorical construction known as the Karoubi envelope/Cauchy completion/splitting idempotents which has been studied within the context of Categorical Quantum Mechanics in [26, 42, 76] and in more general theories in [33, 72, 56, 67, 41].

Intuitively, the Karoubi envelope \mathcal{K} of a process theory, \mathcal{P} , can be thought of as the process theory which contains every possible decohered subtheory \mathcal{P}^{dec} in a minimal way. That is, for any choice of decoherence processes dec , it is the case that $\mathcal{P}^{\text{dec}} \subseteq \mathcal{K}[\mathcal{P}]$. More formally it is defined as follows:

Definition 5 (Karoubi envelope, \mathcal{K}). Systems labeled in $\mathcal{K}[\mathcal{P}]$ are labeled by arbitrary pairs:

$$\left| \left(A, \diamond_A^A \right) \right| \quad (73)$$

where A is a system in \mathcal{P} and \diamond is an idempotent on A . We introduce the shorthand notation for systems:

$$\mathbb{A}_\diamond := \left(A, \diamond_A^A \right) \quad (74)$$

and, noting that there can be multiple different idempotents per system we will distinguish them by their colour, hence we can have systems such as \mathbb{A}_\diamond and \mathbb{A}_\bullet which differ only by their associated idempotent.

We can then note that these compose via:

$$\left| \left(A, \diamond_A^A \right) \right| \left| \left(B, \diamond_B^B \right) \right| := \left| \left(AB, \diamond_A^A \diamond_B^B \right) \right| \quad (75)$$

Processes in $\mathcal{K}[\mathcal{P}]$ from \mathbb{A}_\diamond to \mathbb{B}_\bullet are a subset of the processes in \mathcal{P} from A to B , namely, those satisfying:

$$\begin{array}{c} |B \\ \boxed{f} \\ |A \end{array} = \begin{array}{c} |B \\ \blacklozenge \\ \boxed{f} \\ |A \\ \diamond_A^A \end{array} \quad (76)$$

Processes in $\mathcal{K}[\mathcal{P}]$ then compose exactly as their associated processes would in \mathcal{P} .

The connection to decoherence is clear – the processes in $\mathcal{K}[\mathcal{P}]$ from \mathbb{A}_\diamond to \mathbb{B}_\bullet correspond to decohered processes where system A has been decohered by \diamond and system B by \blacklozenge . This is in this way that the process theory contains every possible decohered system from \mathcal{P} . Note, in particular, that this means that it contains \mathcal{P} itself which is simply corresponds to restricting to systems of the form $(A, |_A)$.

Typically it will be the case that $\mathcal{P} \subsetneq \mathcal{K}[\mathcal{P}]$ and so we cannot directly use the Karoubi envelope to define a partitioning of \mathcal{P} . It is however relatively simple to see how one can modify \mathcal{P} such that we can view $\mathcal{K}[\mathcal{P}]$ as a set of decohered processes and leverage this to define a free partition and so on to resource theories. The way we do this is we extend \mathcal{P} to give systems an additional label, namely, a decoherence process.

Definition 6 (Decoherence labeled theory, \mathcal{D}). Systems in $\mathcal{D}[\mathcal{P}]$ are exactly the same as those in $\mathcal{K}[\mathcal{P}]$, that is, they are labeled by arbitrary pairs:

$$\left| \left(A, \diamond_A^A \right) \right| \quad (77)$$

where A is a system in \mathcal{P} and \diamond is an idempotent on A . Composition similarly is the same as in $\mathcal{K}[\mathcal{P}]$ that is:

$$\left| \left(A, \diamond_A^A \right) \right| \left| \left(B, \diamond_B^B \right) \right| := \left| \left(AB, \diamond_A^A \diamond_B^B \right) \right| \quad (78)$$

The difference between $\mathcal{K}[\mathcal{P}]$ and $\mathcal{D}[\mathcal{P}]$ comes when we look at the processes. Specifically, processes in $\mathcal{D}[\mathcal{P}]$ from \mathbb{A}_\diamond to \mathbb{B}_\bullet are the full set of processes in \mathcal{P} from A to B – the decoherence map is simply a label and does not have an impact on the set of possible processes. These then compose exactly as their associated processes would in \mathcal{P} .

That is, we take the same systems as in the Karoubi envelope but we do not restrict the processes to being the decohered processes, we allow them to be arbitrary processes from \mathcal{P} . It then immediately follows that $\mathcal{K}[\mathcal{P}] \subseteq \mathcal{D}[\mathcal{P}]$ and hence that this defines a partitioning of $\mathcal{D}[\mathcal{P}]$. This, for the same reasons as discussed in Sec. 3, will not lead to interesting resource theories of states and so we define a broader class of processes for the free subtheory.

Following our earlier discussions it is natural to define the analogue of cDIO, however, note that the assumption about how systems compose in $\mathcal{D}[\mathcal{P}]$ (as described by Eq. (78)) and similarly in $\mathcal{K}[\mathcal{P}]$ (as described by Eq. (75)) means that we are in the situation described by Ex. 4 and so, as discussed just after Lem. 4 it will be the case here that cDIO = DIO. We will therefore define the following as a potential set of free processes:

Definition 7. Given a decoherence labeled process theory $\mathcal{D}[\mathcal{P}]$ we can define the subtheory DIO(= cDIO) as:

$$\text{DIO} := \left\{ \begin{array}{c} \begin{array}{c} |B \\ \boxed{g} \\ |A_\diamond \end{array} \in \mathcal{D}[\mathcal{P}] \quad \left| \quad \begin{array}{c} |B \\ \boxed{g} \\ |A_\diamond \\ \diamond \\ |A_\diamond \end{array} = \begin{array}{c} |B \\ \boxed{g} \\ |A_\diamond \end{array} \end{array} \right\} \quad (79)$$

This defines a partitioned process theory (DIO, $\mathcal{D}[\mathcal{P}]$), which can then be used to construct various resource theories of coherence following the methods outlined in App. A.

The resource theories which arise from this partitioned process theory will be much more elaborate than those considered so far in the literature. This is because they do not pick a particular decoherence mechanism for each system, but allow one to consider every possible decoherence mechanism for each system. Even in the simplest case of constructing a resource theory of states for this partitioned process theory, we find a substantially richer structure than has so far been considered in the literature – developing a deeper understanding of this structure will be left for future work.

Note that it is possible to go beyond the assumption that global decoherence is simply the product of local decoherences which led to the fact that cDIO = DIO. To do so would involve redefining how systems compose within $\mathcal{K}[\mathcal{P}]$ and $\mathcal{D}[\mathcal{P}]$. That is, we would take some new composition rule:

$$\left(\begin{array}{c} |A \\ \diamond \\ |A \end{array} \right) \left(\begin{array}{c} |B \\ \diamond \\ |B \end{array} \right) := \left(\begin{array}{c} |A \quad |B \\ \diamond \\ |A \quad |B \end{array} \right) \quad (80)$$

where the global decoherence

$$\begin{array}{c} |A \quad |B \\ \diamond \\ |A \quad |B \end{array} \quad (81)$$

satisfies certain constraints, for example, those discussed in Ex. 5. Again, we leave detailed analysis of this for future work.

Similarly, we can go beyond the assumption of speakable coherence to the more general unspeakable case. We simply must replace decoherence processes in the above constructions with decoherence mechanism – i.e., in the end we will have one system per decoherence mechanism and a commuting condition defining the free processes. Once again, detailed analysis of this will be left for future work.

5 Discussion

In the previous section we illustrated the use of the process theoretic approach for constructing resource theories of coherence. As we explained, this should be seen as a toolbox for constructing these resource theories: if you have a particular situation in mind then the above techniques should inform how to go about constructing the relevant resource theory, with the examples we have given serving as a guide. In constructing a resource theory of coherence there are many decisions that must be made:

- What are the decoherence processes, and, what compositionality constraints do they need to satisfy?
- Is the coherence speakable or unspeakable?
- If the answer is unspeakable coherence, then what are the decoherence mechanisms? e.g. how much and what information should be encoded in the environment?
- Finally, what are the fundamental resources to be considered, are they states, measurements, transformations or something more general?

Once these decisions have been made, then all of the machinery that we've introduced in this work can be employed to construct the appropriate resource theory.

Perhaps the most important lessons to be learnt here, which apply regardless of the above decisions that have to be made, are that:

- One should demand that the free processes preserve the full set of decohered processes rather than just states—this rules out MIO as a suitable choice for free processes.
- Considering composite systems has an influence even on single-system resource theories—showing that we need to move from considering DIO processes to a 'complete' notion which we term cDIO.

These tell us that the largest set of free processes we should consider are the cDIO processes we defined.

Whilst cDIO processes are singled out as the largest set of free processes, if we want to consider unspeakable rather than speakable coherence, it is well motivated to consider a smaller set of free processes, we show the essential difference here comes when one considers the decoherence mechanism rather than just the decoherence process. This is closely related to the translationally invariant operations, TIO, which have widely been considered in the literature, again with a caveat that they should be invariant in a complete sense, we call such processes cTIO processes. We find that this is in really a limiting case of the general construction which allows for us to interpolate between cDIO and cTIO by varying how much information is encoded into the environment system.

We have shown how we can, in an extremely general way, talk about resource theories of coherence for process theories. Interestingly, this makes no reference to tomography, convexity, or any of the usual structure of generalised probabilistic theories, it would therefore be interesting to know if coherence plays a role in fields which are normally well beyond the scope of what we consider in quantum foundations. Moreover, this challenges the standard tools used to study coherence in quantum information theory, no longer are bases playing a prominent role, or any convex or information theoretic structures. This opens the question as to how far we can push this approach, what else can we learn from this perspective?

5.1 Future work

There are plenty of directions for future work stemming from this, we hope this sparks interest in these tools and they are explored more broadly by the community.

Firstly we can consider applications of this formalism:

- On the one hand, there is the question of applying this formalism specifically to quantum theory – to do a more in depth study of the various resource theories that can be constructed, to understand the properties that they have, how these relate to information theoretic tasks, and how to define meaningful monotones to capture the preorder of resources. In particular, i) to explore decoherence maps beyond the traditionally studied completely dephasing maps, ii) to explore resource theories with multiple sources of decoherence, and iii) to use these tools to, in a methodical way, go beyond considering coherence of states to coherence of channels and more general objects.
- On the other hand, we can consider applying this formalism to other physical theories, for example, what can we say about coherence in generalised probabilistic theories [8], in Spekkens toy model [82], or in more exotic process theories such as those described in [32].

Secondly, we can consider further development of the formalism itself:

- In one direction, it would be beneficial to explore the relationship between all of the different resource theories that we can construct using these tools – are there some basic principles which allow us to organise this ‘zoo’ of resource theories? Moreover, given that we have multiple different ways to define the partitioned process theory, then it would be worth exploring the basic mathematical structure of this as a play ground for defining resource theories with multiple different resources. This should have connections to the work of [79, 80]. Indeed, section 4.5 can be viewed as preliminary work in this direction – at least for the case of decoherence.
- In another direction we can consider generalising our notion of decoherence. At times it may be useful to consider decoherence not as being a physical process, but simply as some restriction on what some agent can do. For example viewing the rebit [40] as a ‘decohered’ form of a qubit. Such a picture strongly connects to the work of [12] and could be modelled by defining decoherence directly on the level of the process theory as an idempotent semi-functor. This would avoid the necessity for decoherence to be a physically implementable process and so, for example, this would allow one to view the rebit as being a decohered form of a qubit. More generally, if we do not demand that decoherence is physical then we would recover the notion of a *resource destroying map* [59] our condition that a process is free relative to the resource destroying map would be a complete version of the commutativity condition in [59].

Finally, we hope that this sparks interest in the process-theoretic approach to resource theories, and moreover highlights the importance of considering composition in a much more general sense than is typically considered in the literature. We have seen that such consideration lead to important restrictions on the free processes for resource theories of coherence, and they may have important consequences for many other resource theories.

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A Resource theories

As discussed in the introduction, in Ref. [23] a resource theory, \mathcal{R} , is defined to be formally equivalent to a process theory; however, rather than interpreting systems as being physical systems and processes as their evolution, we instead interpret systems as being resources, and processes as ways in which we can freely transform between resources. Composite systems simply represent having multiple different resources available. Sequential composition of the processes reflects the fact that if you can freely transform one resource into another and then freely transform that into a third, then there is a way to freely transform the first resource into the third resource. A similar story can be told for parallel composition. A resource is then said to be free if there is some way to freely create the resource out of nothing.

To try to give some intuition for this distinction between process theories and resource theories, let us consider two different presentations of an LOCC transformation in the resource theory of entanglement. Specifically, let us consider the transformation of a maximally entangled Bell state into a product state by using local discard and prepare channels. On the one hand we can represent this in the process theory of quantum theory:

$$\begin{array}{c} \begin{array}{c} \text{C}^2 \\ \downarrow \\ \triangle 0 \\ \hline \text{C}^2 \end{array} \quad \begin{array}{c} \text{C}^2 \\ \downarrow \\ \triangle 0 \\ \hline \text{C}^2 \end{array} \\ \hline \triangle \text{Bell} \end{array} \quad (82)$$

where this diagram describes the physical procedure by which the Bell state is prepared, which is then composed with a discard and prepare $|00\rangle$ channel.

On the other hand we could consider the resource theory of LOCC-entanglement where the same procedure would be represented as:

$$\begin{array}{c} |00\rangle \\ \downarrow \\ \boxed{\text{Discard-Prepare } |00\rangle} \\ \downarrow \\ |\text{Bell}\rangle \end{array} \quad (83)$$

where this diagram should be read as saying that, if one had access to the resource of ‘a Bell state’, then there is a free way to transform it under LOCC to the resource ‘a $|00\rangle$ state’.

There is, of course, a much simpler way to obtain the resource ‘a $|00\rangle$ state’ from LOCC operations which is simply to directly prepare it. In our process theory of quantum theory this would be represented by:

$$\begin{array}{c} \text{C}^2 \\ \downarrow \\ \triangle 0 \end{array} \quad \begin{array}{c} \text{C}^2 \\ \downarrow \\ \triangle 0 \end{array} \quad (84)$$

whilst in the resource theory of LOCC-entanglement this would be represented by:

$$\begin{array}{c} |00\rangle \\ \downarrow \\ \boxed{\text{Prepare } |00\rangle} \end{array} \quad (85)$$

the fact that this resource can be constructed in the resource theory from ‘nothing’ means that ‘a $|00\rangle$ state’ is a *free resource*.

In other words, whilst resource theories and process theories may be formally equivalent mathematical structures we should be extremely careful not to conflate them. There is however a strong connection that can be identified between these two perspectives, indeed the main contribution of [23] was identifying the relevant structure that needs to be added to a process theory in order to construct a resource theory from it.

Specifically, given a process theory \mathcal{P} , in order to construct resource theories we must identify a sub-process theory of free processes $\mathcal{P}^{\text{free}}$. The idea being that these are the physical processes which we can freely implement with no cost. For example, in constructing a resource theory of entanglement we may identify the LOCC operations as being free. This subset of free processes should be closed under composition, that is, if we can freely implement f and freely implement g then we should be able to freely do a composite of f and g . Hence, rather than being an arbitrary subset of the processes this subset must be a sub process theory. Such a pair $(\mathcal{P}, \mathcal{P}^{\text{free}})$ is known as a *partitioned process theory*.

Given a partitioned process theory $(\mathcal{P}, \mathcal{P}^{\text{free}})$ they show in [23] that there are various different resource theories which can be constructed. Roughly speaking, this consists of identifying the sorts of processes within \mathcal{P} which you want to consider to be resources. For example, do we want to think of just the states in \mathcal{P} as being resources, or do we want to think of the channels as being resources? There are various constructions of this sort which are discussed in [23], we, for simplicity, will just discuss a few of these here. But it should be noted that this is not exhaustive, and neither was the original work. The precise construction that should be used will depend on exactly what one chooses to designate as a resource. In this sense, the partitioned process theory is really the more fundamental structure, and the particular resource theory that one constructs is situation dependent.

Example 7 (Resource theory of states). Firstly let us construct a resource theory of *states*. These are the most commonly considered resource theories within the literature. In such resource theories, the systems in the resource theory \mathcal{R} correspond to all of the states in \mathcal{P} :

$$\left\{ \left| s^A \right. \right\} \cong \left\{ \left\langle \begin{array}{c} \downarrow \\ s \end{array} \right| \begin{array}{c} \uparrow \\ A \end{array} \right\rangle \in \mathcal{P} \right\} \quad (86)$$

where parallel composition of systems in \mathcal{R} is given by parallel composition of the associated states in \mathcal{P} , that is:

$$\left| s^A \right. \left| r^B \right. \sim \left\langle \begin{array}{c} \downarrow \\ s \end{array} \right| \begin{array}{c} \uparrow \\ A \end{array} \right\rangle \left\langle \begin{array}{c} \downarrow \\ r \end{array} \right| \begin{array}{c} \uparrow \\ B \end{array} \right\rangle \quad (87)$$

whilst the transformations correspond to those in $\mathcal{P}^{\text{free}}$:

$$\left\{ \left[\begin{array}{c} \left| r^B \right. \\ \boxed{f} \\ \left| s^A \right. \end{array} \right] \right\} \cong \left\{ \left[\begin{array}{c} \left| B \right. \\ \boxed{f} \\ \left| A \right. \end{array} \right] \in \mathcal{P}^{\text{free}} \left| \left[\begin{array}{c} \left| B \right. \\ \boxed{f} \\ \left| A \right. \\ \left\langle \begin{array}{c} \downarrow \\ s \end{array} \right| \end{array} \right] = \left[\begin{array}{c} \left| B \right. \\ \left\langle \begin{array}{c} \downarrow \\ r \end{array} \right| \end{array} \right] \right\} \quad (88)$$

that is, we can freely transform s^A into r^B if there is some free process with input A and output B that, in particular, maps the state s of A to the state r of B .

Composition of processes in \mathcal{R} is given by the composition of processes in \mathcal{P} , in particular, sequential composition is given by:

$$\begin{array}{c} |t^C \\ \boxed{g} \\ |r^B \\ \boxed{f} \\ |s^A \end{array} \sim \begin{array}{c} |C \\ \boxed{g} \\ |B \\ \boxed{f} \\ |A \end{array} \quad (89)$$

where f transforms s into r and then g transforms r into t in the process theory \mathcal{P} . Parallel composition is given by:

$$\begin{array}{c} |r^B \\ \boxed{f} \\ |s^A \end{array} \quad \begin{array}{c} |u^D \\ \boxed{g} \\ |t^C \end{array} \sim \begin{array}{c} |B \\ \boxed{f} \\ |A \end{array} \quad \begin{array}{c} |D \\ \boxed{g} \\ |C \end{array} \quad (90)$$

where f transforms s into r and g transforms t into u within the process theory \mathcal{P} .

The free resources are then simply the resources which can be constructed from nothing:

$$\left\{ \left| s^A \right| \left| \begin{array}{c} |s^A \\ \boxed{s} \end{array} \right. \in \mathcal{R} \right\} \quad (91)$$

which is clearly nothing but the set of states in the free partition as:

$$\left\{ \left| \begin{array}{c} |s^A \\ \boxed{s} \end{array} \right. \right\} \cong \left\{ \left| \begin{array}{c} |A \\ \triangle \\ |s \end{array} \right. \in \mathcal{P}^{\text{free}} \right\} \quad (92)$$

Example 8 (Resource theory of effects). Recently there has been interested in resource theories of effects [65, 37]. We can capture that in our formalism as follows. Define the resources to simply be the effects in \mathcal{P} :

$$\left\{ \left| e_A \right. \right\} \cong \left\{ \left| \begin{array}{c} \triangle e \\ |A \end{array} \right. \in \mathcal{P} \right\} \quad (93)$$

where parallel composition of systems in \mathcal{R} is given by parallel composition of the associated effects in \mathcal{P} , that is:

$$\left| e_A \right. \left| e'_B \right. \sim \begin{array}{c} \triangle e \\ |A \end{array} \begin{array}{c} \triangle e' \\ |B \end{array} \quad (94)$$

transformations in \mathcal{R} then correspond to processes in $\mathcal{P}^{\text{free}}$ transforming effects into effects by pre-composition:

$$\left\{ \left| \begin{array}{c} |e_A \\ \boxed{f} \\ |e'_B \end{array} \right. \right\} \cong \left\{ \left| \begin{array}{c} |B \\ \boxed{f} \\ |A \end{array} \right. \in \mathcal{P}^{\text{free}} \left| \begin{array}{c} \triangle e' \\ |B \\ \boxed{f} \\ |A \end{array} \right. = \left| \begin{array}{c} \triangle e \\ |A \end{array} \right. \right\} \quad (95)$$

Composition of processes in \mathcal{R} is given by the relevant composition of processes in \mathcal{P} .

Similarly to the case of states, it is easy to show that the free resources in \mathcal{R} will be in one-to-one correspondence with the effects in $\mathcal{P}^{\text{free}}$.

Example 10 (Resource theory of processes). Finally, let us construct a resource theory of general processes. This is less commonly considered in the literature than resource theories of states, the ability for this general compositional framework to so naturally handle this situation is one of its main advantages, there are some notable exceptions however such as [58, 60, 36]. The first difference is that now we label the systems in the resource theory \mathcal{R} by arbitrary processes in \mathcal{P} , that is:

$$\left\{ \left| f_A^B \right. \right\} \cong \left\{ \left[\begin{array}{c} \text{---} B \\ \boxed{f} \\ \text{---} A \end{array} \right] \in \mathcal{P} \right\} \quad (100)$$

with parallel composition of systems being given by parallel composition of the associated processes. Transformations in the resource theory are then given by ‘free 1-combs’, $\xi = (\xi_1 \in \mathcal{P}^{\text{free}}, \xi_2 \in \mathcal{P}^{\text{free}})$, that is:

$$\left\{ \left[\begin{array}{c} \text{---} D \\ \boxed{g_C^D} \\ \boxed{\xi} \\ \text{---} f_A^B \end{array} \right] \right\} \cong \left\{ \left(\left[\begin{array}{c} \text{---} D \\ \boxed{\xi_2} \\ \text{---} B \\ \text{---} E \\ \text{---} A \\ \boxed{\xi_1} \\ \text{---} C \end{array} \right] \in \mathcal{P}^{\text{free}} \mid \left[\begin{array}{c} \text{---} D \\ \boxed{\xi_2} \\ \text{---} B \\ \boxed{f} \\ \text{---} A \\ \boxed{\xi_1} \\ \text{---} C \end{array} \right] = \left[\begin{array}{c} \text{---} D \\ \boxed{g} \\ \text{---} C \end{array} \right] \right\} \quad (101)$$

We can then define sequential composition as:

$$\left[\begin{array}{c} \text{---} h_E^F \\ \boxed{\eta} \\ \text{---} g_C^D \\ \boxed{\xi} \\ \text{---} f_A^B \end{array} \right] \sim \left[\begin{array}{c} \text{---} F \\ \boxed{\eta_2} \\ \text{---} D \\ \boxed{\xi_2} \\ \text{---} B \\ \text{---} X \\ \text{---} Y \\ \text{---} A \\ \boxed{\xi_1} \\ \text{---} C \\ \boxed{\eta_1} \\ \text{---} E \end{array} \right] \quad (102)$$

where ξ maps f to g and then η maps g to h within \mathcal{P} . Parallel composition is defined as:

$$\left[\begin{array}{c} \text{---} g_C^D \\ \boxed{\xi} \\ \text{---} f_A^B \end{array} \right] \left[\begin{array}{c} \text{---} h_E^F \\ \boxed{\eta} \\ \text{---} g_C^D \\ \boxed{\xi} \\ \text{---} f_A^B \end{array} \right] \sim \left[\begin{array}{c} \text{---} D \\ \boxed{\xi_2} \\ \text{---} B \\ \text{---} F \\ \text{---} X \\ \text{---} Y \\ \text{---} A \\ \boxed{\xi_1} \\ \text{---} C \end{array} \right] \left[\begin{array}{c} \text{---} H \\ \boxed{\eta_2} \\ \text{---} H \\ \text{---} X \\ \text{---} Y \\ \text{---} E \\ \boxed{\eta_1} \\ \text{---} G \end{array} \right] \quad (103)$$

where ξ maps f to g and η maps h to i within \mathcal{P} .

By the same logic as in Ex. 9 it can be shown that the free resources are in one-to-one correspondence with the set of processes in $\mathcal{P}^{\text{free}}$.

Whilst this last example may seem very general, there is actually a substantial limitation to this approach. Namely, that it specifies that all of the resources must be processed in parallel. That is, there is no way to take two processes viewed as resources and to compose them in sequence to obtain a new resource.

Example 11 (Resource theory of sets of processes). A resource in this theory will be labeled by an indexed set of processes in \mathcal{P} :

$$\left| \{f_i \in \mathcal{P}\}_{i \in I} \right. \quad (104)$$

and parallel composition of these resources will be given by the disjoint union of the sets:

$$\{f_i\}_{i \in I} \otimes \{g_j\}_{j \in J} := \{f_i\}_{i \in I} \sqcup \{g_j\}_{j \in J} \quad (105)$$

For a full formalisation of this see [23], intuitively, however a free transformation $\xi : \{f_i\}_{i \in I} \rightarrow \{g_j\}_{j \in J}$ is given by:

- a disjoint partitioning of $I = \bigsqcup_{j \in J} I_j$
- for each $j \in J$ a circuit ξ_j constructed out of processes in $\mathcal{P}^{\text{free}}$ such that
 - there is a ‘hole’ in the circuit for each f_i where $i \in I_j$, and
 - when all of the f_i are ‘plugged into these holes’ the resulting process is g_j

One can then show that the free resources in such a resource theory are simply given by:

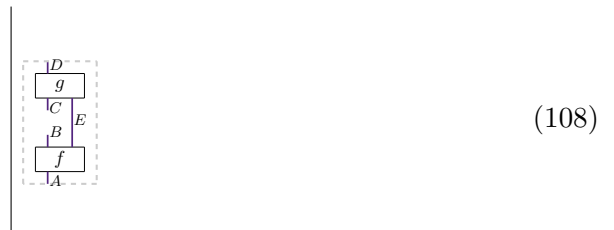
$$\left| \{f_i \in \mathcal{P}^{\text{free}}\}_{i \in I} \right. \quad (106)$$

So far we have seen that we define free processings in terms of higher-order processes, such as the combs in 10 and the circuits with holes in 11, we can also consider resource theories in which the resources themselves are higher-order transformations, for example:

Example 12 (Resource theory of 2-combs). Let us consider resources to be combs of the shape:



that is, we will label the wires of the resource theory by these:



These have been far less widely studied in the literature, however [28] can be viewed as a special case of this resource theory in which E is assumed to be trivial, f is assumed to be a state and g is assumed to be a measurement.

Parallel composition of resources is given by:

(109)

The free transformations, ξ , within the resource theory are in one-to-one correspondence with particular diagrams constructed out of processes in $\mathcal{P}^{\text{free}}$, for example, diagrams of the forms:

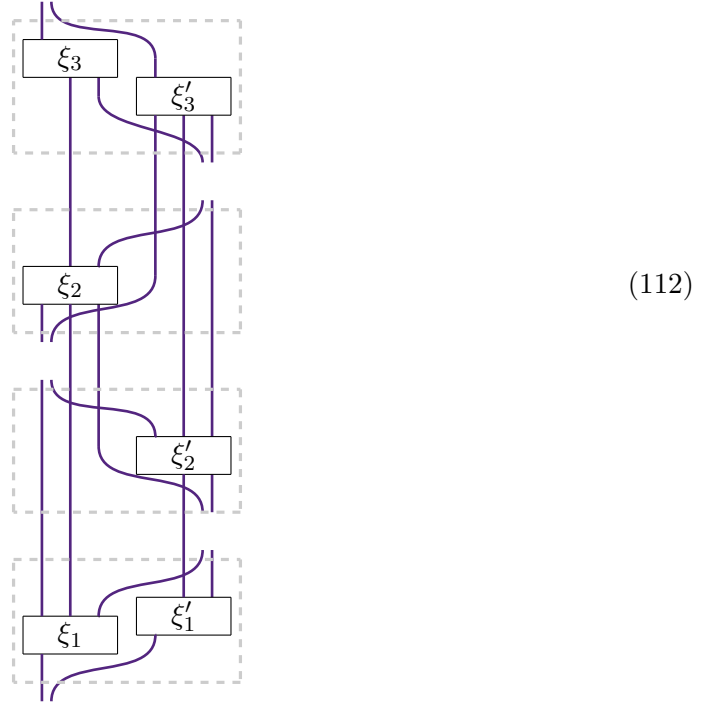
(110)

where $\xi_i \in \mathcal{P}^{\text{free}}$, correspond to transformations for the resource theory \mathcal{R} .

These, for example, can be composed in sequence as:

(111)

and in parallel as:



Finally, the free resources are those that can be freely constructed from nothing, that is, resources:



in which $f, g \in \mathcal{P}^{\text{free}}$.

This is just a few examples of the sorts of resource theories that can be constructed from a partitioned process theory. It is clearly impossible to make an exhaustive list of examples as there is an infinite number of different sorts of things that one may want to consider as resources depending on the physical situation on hand. However, in future work it would be interesting to explore the scope of such resource theories and to develop a formalism which handles the construction of any resource theory from a partitioned process theory in a uniform and principled way.

As mentioned in the introduction, in the study of resource theories we are often not interested in the details of how we transform one resource into another, but simply whether or not there exists some way to freely transform one resource into another. That is, we often work simply with the preorder, R , that these resource theories, \mathcal{R} , define. For example, in the case of the resource theory of states, we work with the preorder defined by:

$$s^A \succ_R r^B \iff \exists (f : s^A \rightarrow r^B) \in \mathcal{R} \quad (114)$$

recall that this means that $f \in \mathcal{P}^{\text{free}}$ within the partitioned process theory. Similarly, in the resource theory of processes we work with the preorder defined by:

$$f_A^B \succ_R g_C^D \iff \exists (\xi : f_A^B \rightarrow g_C^D) \in \mathcal{R} \quad (115)$$

recall also that this means that ξ can be viewed a ‘comb’ constructed out of processes in $\mathcal{P}^{\text{free}}$. Essentially, in order to move from the resource theory, \mathcal{R} , to the preorder, R , we simply forget all of details of the resource theory except for whether the set of transformations from one resource to another is the empty set or not.

Often we are not able to efficiently characterise this preorder, in which case the study of resource theories often revolves around the construction of *monotones*. That is, a function $\alpha : R \rightarrow \mathbb{R}$ such that:

$$s \succ_R r \implies \alpha(s) \geq \alpha(r) \quad (116)$$

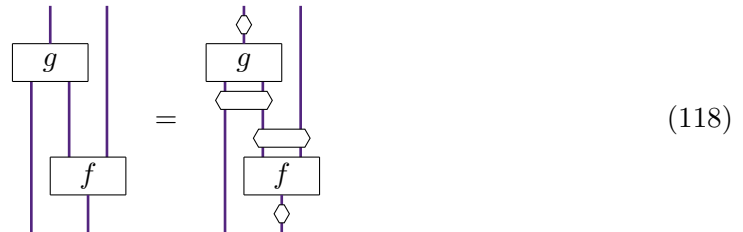
These usual aim for such monotones is to find out some information about the preorder, in a way that is simple to compute and has some physical or operational significance. For this work we simply aim for an ‘in principle’ characterisation of \mathcal{R} and hence R rather than trying to find efficient characterisations of this structure via monotones.

B Proof of lemma 1

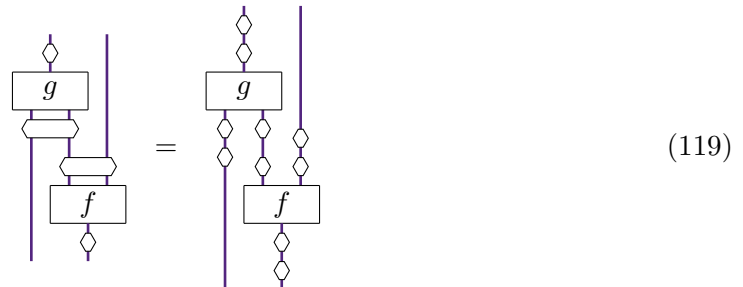
Proof. To see that this enforces that the composition of decohered is itself a decohered process it suffices to consider a pair of decohered processes f and g composed as:



as all other ways of composing processes can be seen as special cases of this. Then, as they are decohered we can write that:



Then using constraint (24) together with idempotence of decoherence processes we find that:



Then, using the fact that f and g are decohered for a second time, and a simple diagram-

matic rewriting, we obtain:

Diagrammatic equation (120) showing the commutation of decoherence diamonds with process boxes g and f . The left side shows a sequence of three vertical lines. The top line has two diamonds, followed by a box g , then two more diamonds, then a box f , and finally two more diamonds. The right side shows the same sequence but with the diamonds moved: the top line has one diamond, then box g , then one diamond; the middle line has one diamond, then box f , then one diamond; and the bottom line has two diamonds. An equals sign is between the two sides.

Finally, using our constraint on composition of decoherences, eq. (24), we can write that:

Diagrammatic equation (121) showing the composition of two decoherence processes. The left side shows two vertical lines. The top line has a diamond, then box g , then a diamond. The bottom line has a diamond, then box f , then a diamond. The right side shows two vertical lines. The top line has a diamond, then box g , then a diamond, and a horizontal bar above it. The bottom line has a diamond, then box f , then a diamond, and a horizontal bar below it. An equals sign is between the two sides.

From which we can conclude that this composite of f and g is also a decohered process. \square

C Proof of lemma 2

Proof. Firstly consider the RHS of the first constraint, eq. (25), and note that, by applying the second constraint (eq. (26)) multiple times, we can write that:

Diagrammatic equation (122) showing the simplification of a stack of decoherence diamonds. The left side shows four horizontal bars stacked vertically on two vertical lines. The middle side shows two horizontal bars on two vertical lines, with a dashed box around them. The right side shows two vertical lines, each with a diamond, and a third vertical line with a diamond. An equals sign is between the first and second sides, and another equals sign is between the second and third sides.

Equating the left hand side using our first constraint (eq. (25)) and using idempotence of decoherence processes on the right hand side then gives us:

Diagrammatic equation (123) showing the idempotence of decoherence processes. The left side shows two horizontal bars stacked vertically on two vertical lines. The right side shows three vertical lines, each with a diamond. An equals sign is between the two sides.

Now, consider the composite of two decohered processes f and g , the fact that they are decohered means that we can write:

Diagrammatic equation (124) showing the decomposition of a composite process into a decohered process. The left side shows two vertical lines. The top line has box g , and the bottom line has box f . The right side shows two vertical lines. The top line has a diamond, then box g , then a diamond, then a horizontal bar, then a diamond. The bottom line has a diamond, then box f , then a diamond, then a horizontal bar, then a diamond. An equals sign is between the two sides.

Now, using the result we just derived (123) we can write that

$$(125)$$

Using idempotence and the above result, that is (123), a second time we obtain:

$$(126)$$

The second constraint (i.e., eq. (26)) then gives us that:

$$(127)$$

Finally, using the fact that f and g are decohered means that we can conclude that:

$$(128)$$

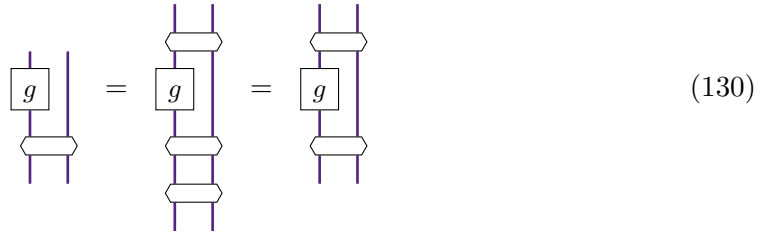
Hence, the composite of f and g is itself a decohered process. \square

D Proof of lemma 4

To see that def. 3 implies def. 4 is simple. Firstly note that $\text{dec} \subseteq \mathcal{P}^{\text{dec}}$ as, due to idempotency of decoherence processes we have that:

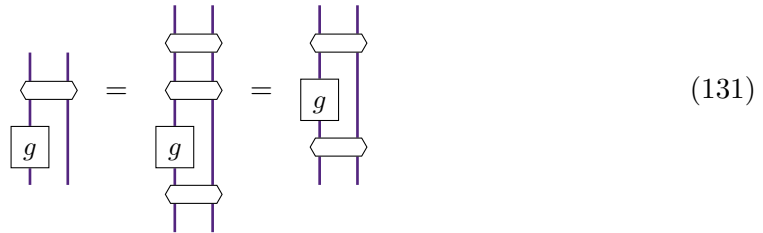
$$(129)$$

therefore, under the assumption of def. 3, we have that the composite of g with this decoherence process must itself be a decohered process, that is:




$$(130)$$

where the first equality is simply the definition of a decohered process, and the second is from idempotency. Similarly, we have:



$$(131)$$

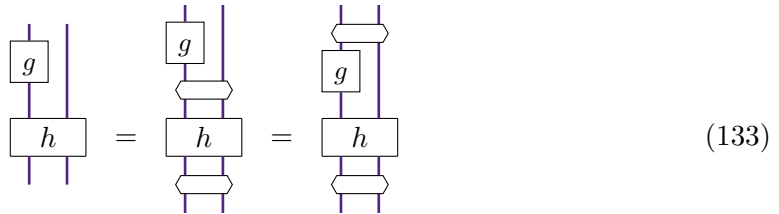
combining these we obtain



$$(132)$$

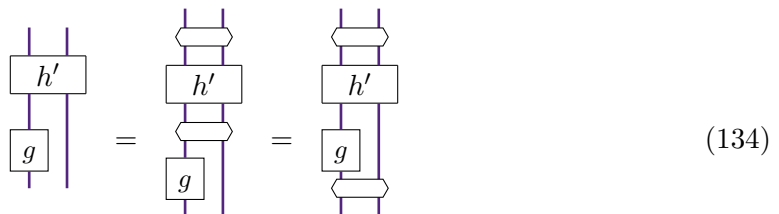
as required to satisfy def. 4.

Now, consider the converse direction. Assume we have some decohered process h , then



$$(133)$$

where the first equality is from the fact that h is decohered and the second from the assumption of def. 4. This means that the composite process is indeed a decohered process as we required. Similarly, given a decohered process h' we can show that:



$$(134)$$

and hence, the minimal constraint on free processes is satisfied.