

the same measurement, but then their corresponding probabilities must sum at most to 1. The second condition follows from the fact that if d vectors are mutually orthogonal in \mathbb{C}^d , they form a complete basis, and then their corresponding probabilities must exactly sum to one. Note that the first condition implies in particular that any two vectors $|v_1\rangle$ and $|v_2\rangle$ in \mathcal{S} that are orthogonal cannot both be assigned the value 1 by f .

We call any assignment $f : \mathcal{S} \rightarrow \{0, 1\}$ satisfying the above two conditions, a $\{0, 1\}$ -coloring of \mathcal{S} . The Kochen-Specker theorem states that if $d \geq 3$, there exist sets of vectors that are not $\{0, 1\}$ -colorable, thus establishing the impossibility of a non-contextual deterministic outcome assignment. We call such $\{0, 1\}$ -uncolorable sets, KS sets. In their original proof, Kochen and Specker describe a set \mathcal{S} of 117 vectors in \mathbb{C}^d dimension $d = 3$ [1]. The minimal KS set contains 18 vectors in dimension $d = 4$ [18, 20].

In this paper, we identify within KS sets interesting subsets which we term 01-gadgets. Such 01-gadgets are $\{0, 1\}$ -colorable and thus do not represent by themselves KS sets. However, they do not admit arbitrary $\{0, 1\}$ -coloring: in any $\{0, 1\}$ -coloring of a 01-gadget, there exist two non-orthogonal vectors $|v_1\rangle$ and $|v_2\rangle$ that cannot both be assigned the color 1. We show that such 01-gadgets form the essence of the KS contradiction, in the sense that every KS set contains a 01-gadget and from every 01-gadget one can construct a KS set.

Besides being useful in the construction of KS sets, we show that 01-gadgets also form a fundamental primitive in constructing statistical KS arguments à la Clifton [17] and state-independent non-contextuality inequalities as introduced in [25]. Moreover, we show that an “extended” Kochen-Specker theorem considered by Pitowsky [22] and Abbott et al. [2, 3] can be easily proven using an extension of the notion of 01-gadgets. We give simple constructive proofs of these different results.

Certain 01-gadgets have already been studied previously in the literature, as they possess other interesting properties. In particular, 01-gadgets were also used in [15] to show that the problem of checking whether certain families of graphs (which represent natural candidates for KS sets) are $\{0, 1\}$ -colorable is NP-complete, a result which we refine in the present paper. Specific 01-gadgets have already been studied in the literature, for instance as ‘definite prediction sets’ in [21] and recently as ‘true-implies-false sets’ in [34] where also minimal constructions in several dimensions were explored. A first method to produce different 01-gadgets was also shown in [32].

This paper is organized as follows. In section 2, we introduce some notation and elementary concepts, in particular the representation of KS sets as graphs. In section 3, we define the notion of 01-gadgets and establish their relation to KS sets. In section 4, we give

several constructions of 01-gadgets and associated KS sets. In section 5, we show how 01-gadgets can be used to construct statistical KS arguments. In section 6, we also show a simple constructive proof of the extended Kochen-Specker theorem of Pitowsky [22] and Abbott et al. [3] using a notion of extended 01-gadgets which we introduce. In section 7, we show that 01-gadgets can be used to establish the NP-completeness of $\{0, 1\}$ -coloring of the family of graphs relevant for KS proofs. We finish by a general discussion and conclusion in section 9.

2 Preliminaries

Much of the reasoning involving KS sets is usually carried out using a graph representation of KS sets defined below. We thus start by reminding some basic graph-theoretic definitions.

Graphs. Throughout the paper, we will deal with simple undirected finite graphs G , i.e., finite graphs without loops, multi-edges or directed edges. We denote $V(G)$ the vertices of G and $E(G)$ the edges of G . If two vertices v_1, v_2 are connected by an edge, we say that they are adjacent, and write $v_1 \sim v_2$.

A subgraph H of G (denoted $H < G$) is a graph formed from a subset of vertices and edges of G , i.e., $V(H) \subseteq V(G)$ and $E(H) \subseteq E(G)$. An induced subgraph K of G (denoted $K \triangleleft G$) is a subgraph that includes all the edges in G whose endpoints belong to the vertex subset $V(K) \subseteq V(G)$, i.e., $E(K) \subseteq E(G)$ with $(v_1, v_2) \in E(K)$ iff $(v_1, v_2) \in E(G)$ for all $v_1, v_2 \in V(K)$.

A clique in the graph G is a subset of vertices $Q \subseteq V(G)$ such that every pair of vertices in Q is connected by an edge, i.e., $\forall v_1, v_2 \in Q$ we have $v_1 \sim v_2$. A maximal clique in G is a clique that is not a subset of a larger clique in G . A maximum clique in G is a clique that is of maximum size in G . The clique number $\omega(G)$ of G is the cardinality of a maximum clique in G .

Orthogonality graphs. The use of graphs in the context of the KS theorem comes from the fact that it is convenient to represent the orthogonality relations in a KS set \mathcal{S} by a graph $G_{\mathcal{S}}$, known as its orthogonality graph [7, 6]. In such a graph, each vector $|v_i\rangle$ in \mathcal{S} is represented by a vertex v_i of $G_{\mathcal{S}}$ and two vertices v_1, v_2 of $G_{\mathcal{S}}$ are connected by an edge if the associated vectors $|v_1\rangle, |v_2\rangle$ are orthogonal, i.e. $v_1 \sim v_2$ if $\langle v_1 | v_2 \rangle = 0$ (for instance the graph in Fig. 1 is the orthogonality graph of the set of vectors given by eq. (5)).

It follows that in an orthogonality graph $G_{\mathcal{S}}$, a clique corresponds to a set of mutually orthogonal vectors in \mathcal{S} . If $\mathcal{S} \subset \mathbb{C}^d$ contains a basis set of d orthogonal vectors, then the maximum clique in $G_{\mathcal{S}}$ is of size $\omega(G_{\mathcal{S}}) = d$.

Coloring of graphs. The problem of $\{0,1\}$ -coloring \mathcal{S} thus translates into the problem of coloring the vertices of its orthogonality graph $G_{\mathcal{S}}$ such that vertices connected by an edge cannot both be assigned the color 1 and maximum cliques have exactly one vertex of color 1. Formally, we say that an arbitrary graph G is $\{0,1\}$ -colorable if there exists an assignment $f : V(G) \rightarrow \{0,1\}$ such that

- $\sum_{v \in Q} f(v) \leq 1$ for every clique $Q \subset V(G)$;
- $\sum_{v \in Q_{\max}} f(v) = 1$ for every maximum clique $Q_{\max} \subset V(G)$. (4)

The KS theorem is then equivalent to the statement that there exist for any $d \geq 3$, finite sets of vectors $\mathcal{S} \subset \mathbb{C}^d$ (the KS sets) such that their orthogonality graph $G_{\mathcal{S}}$ is not $\{0,1\}$ -colorable. Deciding if a given graph G admits a $\{0,1\}$ -coloring is NP-complete [15]. Note that any graph G that is not $\{0,1\}$ -colorable must contain at least two cliques of maximum size $\omega(G)$. Indeed, if a graph G contains a single clique of maximum size it always admits a $\{0,1\}$ -coloring consisting in assigning the value 0 to all its vertices, except for one vertex in the maximum clique that is assigned the value 1.

Orthogonal representations. For a given graph G , an orthogonal representation \mathcal{S} of G in dimension d is a set of non-zero vectors $\mathcal{S} = \{|v_i\rangle\}$ in \mathbb{C}^d obeying the orthogonality conditions imposed by the edges of the graph, i.e., $v_1 \sim v_2 \Rightarrow \langle v_1 | v_2 \rangle = 0$ [28]. We denote by $d(G)$ the minimum dimension of an orthogonal representation of G and we say that G has dimension $d(G)$. Obviously, $d(G) \geq \omega(G)$. A *faithful* orthogonal representation of G is given by a set of vectors $\mathcal{S} = \{|v_i\rangle\}$ that in addition obey the condition that non-adjacent vertices are assigned non-orthogonal vectors, i.e., $v_1 \sim v_2 \Leftrightarrow \langle v_1 | v_2 \rangle = 0$ and that distinct vertices are assigned different vectors, i.e., $v_1 \neq v_2 \Leftrightarrow |v_1\rangle \neq |v_2\rangle$. We denote by $d^*(G)$ the minimum dimension of such a faithful orthogonal representation of G and we say that G has faithful dimension $d^*(G)$.

Given a graph G of dimension $d(G)$, the orthogonality graph $G_{\mathcal{S}}$ of the minimal orthogonal representation \mathcal{S} of G has faithful dimension $d^*(G_{\mathcal{S}}) = d(G)$. The graph $G_{\mathcal{S}}$ can be seen as obtained from G by adding edges (between vertices that are non-adjacent in G , but corresponding to vectors in \mathcal{S} that are nevertheless orthogonal) and by identifying certain vertices (those that correspond to identical vectors in \mathcal{S}). We say that $G_{\mathcal{S}}$ is the faithful version of G .

KS graphs. While the non- $\{0,1\}$ -colorability of a set \mathcal{S} translates into the non- $\{0,1\}$ -colorability of its orthogonality graph $G_{\mathcal{S}}$, the non- $\{0,1\}$ -colorability of an arbitrary graph G translates into the non- $\{0,1\}$ -colorability of one of its orthogonal representations

only if this representation has the minimal dimension $d(G) = \omega(G)$. Indeed, it is only under this condition that the requirement that $\sum_{v \in Q_{\max}} f(v) = 1$ in the definition of the $\{0,1\}$ -coloring of the graph G gives rise to the corresponding requirement that $\sum_{v \in Q_{\max}} f(|v\rangle) = 1$ for its orthogonal representation (if the dimension d is larger than $\omega(G) = |Q_{\max}|$, the $|Q_{\max}| < d$ mutually orthogonal vectors $\{|v\rangle : v \in Q_{\max}\}$ in \mathbb{C}^d do not form a basis).

If a graph G is not $\{0,1\}$ -colorable and has dimension $d(G) = \omega(G)$, it thus follows that its minimal orthogonal representation \mathcal{S} forms a KS set. If in addition $d^*(G) = \omega(G)$, we say that G is a KS graph (this last condition can always be obtained by considering the faithful version of G , i.e., the orthogonality graph $G_{\mathcal{S}}$ of its minimal orthogonal representation \mathcal{S}).

The problem of finding KS sets can thus be reduced to the problem of finding KS graphs. But as we have noticed above, deciding if a graph is $\{0,1\}$ -colorable is NP-complete. In addition, while finding an orthogonal representation for a given graph can be expressed as finding a solution to a system of polynomial equations, efficient numerical methods for finding such representations are still lacking. Thus, finding KS sets in arbitrary dimensions is a difficult problem towards which a huge amount of effort has been expended [21]. In particular, “records” of minimal Kochen-Specker systems in different dimensions have been studied [18], the minimal KS system in dimension four is the 18-vector system due to Cabello et al. [18, 20] while lower bounds on the size of minimal KS systems in other dimensions have also been established.

3 01-gadgets and the Kochen-Specker theorem

We now introduce the notion of 01-gadgets that play a crucial role in constructions of KS sets.

Definition 1. A 01-gadget in dimension d is a $\{0,1\}$ -colorable set $\mathcal{S}_{gad} \subset \mathbb{C}^d$ of vectors containing two distinguished vectors $|v_1\rangle$ and $|v_2\rangle$ that are non-orthogonal, but for which $f(|v_1\rangle) + f(|v_2\rangle) \leq 1$ in every $\{0,1\}$ -coloring f of \mathcal{S}_{gad} .

In other words, while a 01-gadget \mathcal{S}_{gad} admits a $\{0,1\}$ -coloring, in any such coloring the two distinguished non-orthogonal vertices cannot both be assigned the value 1 (as if they were actually orthogonal). We can give an equivalent, alternative definition of a gadget as a graph.

Definition 2. A 01-gadget in dimension d is a $\{0,1\}$ -colorable graph G_{gad} with faithful dimension $d^*(G_{gad}) = \omega(G_{gad}) = d$ and with two distinguished non-adjacent vertices $v_1 \approx v_2$ such that $f(v_1) + f(v_2) \leq 1$ in every $\{0,1\}$ -coloring f of G_{gad} .

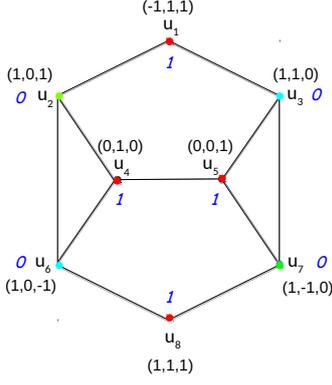


Figure 1: The 8-vertex “Clifton” graph that was used by Kochen and Specker in their construction of the 117 vector KS set. The two distinguished vertices are u_1 and u_8 .

In the following when we refer to a 01-gadget, we freely alternate between the equivalent set or graph definitions.

An example of a 01-gadget in dimension 3 is given by the following set of 8 vectors in \mathbb{C}^3 :

$$\begin{aligned}
 |u_1\rangle &= \frac{1}{\sqrt{3}}(-1, 1, 1), & |u_2\rangle &= \frac{1}{\sqrt{2}}(1, 1, 0), \\
 |u_3\rangle &= \frac{1}{\sqrt{2}}(0, 1, -1), & |u_4\rangle &= (0, 0, 1), \\
 |u_5\rangle &= (1, 0, 0), & |u_6\rangle &= \frac{1}{\sqrt{2}}(1, -1, 0), \\
 |u_7\rangle &= \frac{1}{\sqrt{2}}(0, 1, 1), & |u_8\rangle &= \frac{1}{\sqrt{3}}(1, 1, 1), \quad (5)
 \end{aligned}$$

where the two distinguished vectors are $|v_1\rangle = |u_1\rangle$ and $|v_2\rangle = |u_8\rangle$. Its orthogonality graph is represented in Fig. 1. It is easily seen from this graph representation that the vertices u_1 and u_8 cannot both be assigned the value 1, as this then necessarily leads to the adjacent vertices u_4 and u_5 to be both assigned the value 1, in contradiction with the $\{0, 1\}$ -coloring rules. This graph was identified by Clifton, following work by Stairs [17, 26], and used by him to construct statistical proofs of the Kochen-Specker theorem. We will refer to it as the Clifton gadget G_{Clif} . The Clifton gadget and similar gadgets were termed “definite prediction sets” in [21].

We identify the role played by 01-gadgets in the construction of Kochen-Specker sets via the following theorem.

Theorem 1. *For any Kochen-Specker graph G_{KS} , there exists a subgraph $G_{\text{gad}} < G_{\text{KS}}$ with $\omega(G_{\text{gad}}) = \omega(G_{\text{KS}})$ that is a 01-gadget. Moreover, given a 01-gadget G_{gad} , one can construct a KS graph G_{KS} with $\omega(G_{\text{KS}}) = \omega(G_{\text{gad}})$.*

The demonstration of our theorem is constructive, it allows to build a 01-gadget from a KS graph

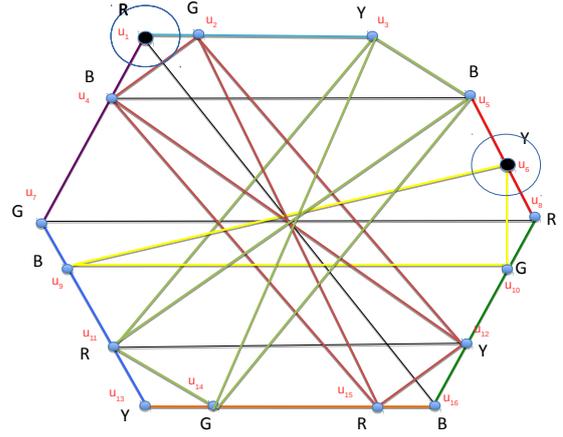


Figure 2: A 16 vertex coloring gadget (also a 101-gadget) that is a subgraph of the 18 vertex Kochen-Specker graph in dimension $d = 4$ found by Cabello et al. [18]. The 9 edge colors denote 9 cliques in the graph, with the maximum clique being of size $\omega(G) = 4$. The distinguished vertices u_1, u_6 are denoted by black circles.

and conversely. The 01-gadget in the original 117-vector proof by Kochen-Specker is the Clifton graph in Fig. 1. A 16-vertex 01-gadget in dimension 4 that is an induced subgraph of the 18-vertex KS graph introduced in [18] is represented in Fig. 2.

Proof. We start by showing the first part of the Theorem: that one can construct a 01-gadget G_{gad} from any KS graph G_{KS} . Given G_{KS} , which by definition is not $\{0, 1\}$ -colorable, we first construct, by deleting vertices one at a time, an induced subgraph G_{crit} that is vertex-critical. By vertex-critical, we mean that (i) G_{crit} is not $\{0, 1\}$ -colorable, but (ii) any subgraph obtained from it by deleting a supplementary vertex does admit a $\{0, 1\}$ -coloring. Observe that in the process of constructing G_{crit} we are able to preserve the maximum clique size, i.e., $\omega(G_{\text{crit}}) = \omega(G_{\text{KS}})$. This is because we are able to delete vertices from all but two maximum cliques, simply because at least two maximum cliques must exist in a graph that is not $\{0, 1\}$ -colorable. Observe also that G_{crit} is itself a KS graph, since the faithful orthogonal representation of G_{KS} in dimension $d = \omega(G_{\text{KS}})$ provides an orthogonal representation of G_{crit} in the same dimension.

We consider three cases: (i) there exists a vertex in G_{crit} that belongs to a single maximum clique, (ii) all vertices in G_{crit} belong to at least two maximum cliques, and there exists a vertex that belong to exactly two maximum cliques; (iii), all vertices in G_{crit} belong to at least three maximum cliques. In the first two cases, which happen to be the case encountered in all known KS graphs, we will be able to prove that the 01-gadget appears as an induced subgraph while

in the third case, the 01-gadget appears as a subgraph that may not necessarily be induced.

In case (i), let v be one of the vertices having the property that it belongs to a single maximum clique. We denote this clique $Q_1 \subset G_S^{\text{crit}}$. Deleting v leads to a graph $G_{\text{crit}} \setminus v$ that is $\{0, 1\}$ -colorable by definition. However, observe that in any coloring f of $G_{\text{crit}} \setminus v$, all the vertices in $Q_1 \setminus v$ are assigned the value 0 by f . This is because, if one of these vertices were assigned value 1, then one could obtain a valid coloring of G_{crit} from f by defining $f(v) = 0$. Choose a vertex $v_1 \in Q_1 \setminus v$ and any other non-adjacent vertex $v_2 \in G_{\text{crit}} \setminus v$. Then $G_{\text{crit}} \setminus v$ is the required 01-gadget with v_1, v_2 playing the role of the distinguished vertices.

In case (ii), let v be one of the vertices having the property that it belongs to exactly two maximum cliques, which we denote $Q_1, Q_2 \subset G_{\text{crit}}$. Again, deleting v leads to a graph $G_{\text{crit}} \setminus v$ that is $\{0, 1\}$ -colorable. However, in any coloring f of $G_{\text{crit}} \setminus v$, it cannot be that a value $f(v_1) = 1$ and a value $f(v_2) = 1$ are simultaneously assigned to a vertex $v_1 \in Q_1 \setminus v$ and a vertex $v_2 \in Q_2 \setminus v$. This is again because if there was such a coloring f , then one could obtain a valid coloring for G_{crit} by defining $f(v) = 0$, in contradiction with the criticality of G_{crit} . Choose $v_1 \in Q_1 \setminus v$ and $v_2 \in Q_2 \setminus v$ such that v_1 and v_2 are not adjacent. Two such vertices must exist. Indeed, if all vertices $Q_1 \setminus v$ were adjacent to all vertices of $Q_2 \setminus v$, then the maximum clique size would be strictly greater than $\omega(G_{\text{crit}})$. Therefore, we have that $G_{\text{crit}} \setminus v$ is the required 01-gadget with v_1, v_2 the distinguished vertices.

Finally, we consider the case (iii) where each vertex in G_{crit} belongs to at least three maximum cliques. In this case, we cannot proceed as above where we remove a certain vertex v and pick vertices from two maximal cliques containing v , because we can no longer guarantee that these two vertices cannot simultaneously be assigned the value 1 (we can only guarantee that a certain t -uple of vertices, each one picked from the t maximum cliques to which v belongs, cannot all simultaneously be assigned the value 1, which may be thought of as a generalization of the 01-gadget property to t distinguished vertices in place of two). Instead, we proceed as follows. We start by deleting edges of G_{crit} one at a time, to construct a new graph G'_{crit} that is edge-critical. By edge-critical, we mean, similarly to the construction above, that G'_{crit} is not $\{0, 1\}$ -colorable, but any graph obtained from it by deleting a supplementary edge (and thus also by deleting a supplementary vertex) does admit a $\{0, 1\}$ -coloring. As above, we are able to preserve the maximum clique size in the process, i.e., $\omega(G'_{\text{crit}}) = \omega(G_{\text{crit}}) = \omega(G_{\text{KS}})$, and G'_{crit} is still a non- $\{0, 1\}$ -colorable KS graph.

Case (iii a): If the resulting graph G'_{crit} is as in the cases (i) and (ii) above, we proceed as before to construct a 01-gadget from a graph $G'_{\text{crit}} \setminus v$, with

the caveat that choosing two non-adjacent vertices v_1 and v_2 in G'_{crit} does not necessarily guarantee that they correspond to non-orthogonal vectors in the natural representation induced by the one of G_{KS} . This is because we have been removing edges from G_{crit} to construct G'_{crit} . However, we can always choose two vertices v_1 and v_2 that were non-adjacent in the original graph G_{KS} and that thus correspond to non-orthogonal vectors. Again, this is because otherwise the maximum clique size of G_{KS} would be greater than $\omega(G_{\text{KS}})$. Now, in any $\{0, 1\}$ -coloring of $G'_{\text{crit}} \setminus v$, we cannot have both $f(v_1) = 1$ and $f(v_2) = 1$, so that $G'_{\text{crit}} \setminus v$ forms a subgraph of G that is a 01-gadget. But notice that the $\{0, 1\}$ -colorings of $G_{\text{crit}} \setminus v$ are a subset of the $\{0, 1\}$ -colorings of $G'_{\text{crit}} \setminus v$. So that we cannot have both $f(v_1) = 1$ and $f(v_2) = 1$ in any $\{0, 1\}$ -coloring of $G_{\text{crit}} \setminus v$ as well. So that the 01-gadget is given by $G_{\text{crit}} \setminus v$ in this case with v_1, v_2 the distinguished vertices.

Case(iii b): If the resulting graph G'_{crit} is not as in the cases (i) and (ii) above, we proceed as follows. Let v be an arbitrary vertex of G'_{crit} . By assumption, this vertex belong to at least two maximum cliques Q_1, Q_2 (and actually even at least a third one). Delete all the edges (v, v') from Q_1 where $v' \in Q_1$ to form $G'_{\text{crit}} \setminus E_v(Q_1)$ (where $E_v(Q_1)$ denotes the edges incident on v in Q_1) which is $\{0, 1\}$ -colorable by definition. In any such coloring f , either $f(v) = 0$ or $f(v) = 1$. In the first case, we must necessarily have that $f(v') = 0$ for all $v' \in Q_1 \setminus v$, since otherwise the coloring f would also define a valid coloring for G'_{crit} . In the second case, we have $f(v'') = 0$ for all $v'' \in Q_2 \setminus v$ by definition of a coloring. We thus conclude that it cannot be simultaneously the case that $f(v') = f(v'') = 1$ for $v' \in Q_1 \setminus v$ and $v'' \in Q_2 \setminus v$. Choose $v_1 \in Q_1$ and $v_2 \in Q_2$ non-adjacent in G_{KS} , which is always possible by the same argument as given before. The faithful version of the graph $G'_{\text{crit}} \setminus E_v(Q_1)$ forms the required 01-gadget with v_1, v_2 being the distinguished vertices. Indeed, by the preceding argument, one can restore edges from G_{crit} to the graph $G'_{\text{crit}} \setminus E_v(Q_1)$ to obtain the 01-gadget so long as the graph is $\{0, 1\}$ -colorable, an instance of this is the graph $G'_{\text{crit}} \setminus (v, v_1)$.

We now proceed to prove the second part of the statement. Starting from a gadget graph we give a construction of a KS graph. The construction generalizes the original Kochen-Specker construction of [1] to arbitrary dimensions and arbitrary repeating gadget units. Given G_{gad} , we know that there exists a faithful orthogonal representation $\{|v_i\rangle\}_{i=1}^n$ in a Hilbert space of dimension $d = \omega(G_{\text{gad}})$ with $n = |V(G_{\text{gad}})|$. Let v_1, v_2 denote the distinguished vertices, and let $|v_2^\perp\rangle$ denote a vector orthogonal to $|v_2\rangle$ that lies in the plane $\text{span}(|v_1\rangle, |v_2\rangle)$, spanned by the vectors $|v_1\rangle$ and $|v_2\rangle$, with $\theta = \arccos|\langle v_1 | v_2^\perp \rangle| > 0$ by definition of a 01-gadget. We consider the following cases: (i) $\frac{\pi}{2\theta}$ is rational and can be written as $\frac{p}{q}$ with q an odd integer, (ii) $\frac{\pi}{2\theta}$ is rational and is given by $\frac{p}{q}$ with q an

even integer, or alternatively, $\frac{\pi}{2\theta}$ is irrational.

Case (i): $\frac{\pi}{2\theta}$ is rational and is given by $\frac{p}{q}$ with q an odd integer. Recall that $|v_2^\perp\rangle$ is orthogonal to $|v_2\rangle$ in the plane $\text{span}(|v_1\rangle, |v_2\rangle)$. In the subspace orthogonal to $\text{span}(|v_1\rangle, |v_2\rangle)$, choose a basis consisting of $d-2$ mutually orthogonal vectors $|w_1\rangle, \dots, |w_{d-2}\rangle$. Denoting G'_{gad} as the orthogonality graph of the entire set of these vectors $\{|v_i\rangle\}_{i=1}^n \cup \{|v_2^\perp\rangle, |w_1\rangle, \dots, |w_{d-2}\rangle\}$, we obtain a gadget graph that can be used as a building block in a Kochen-Specker type construction. In particular, the crucial property of G'_{gad} is that in any $\{0,1\}$ -coloring f , $f(v_1) = 1 \Rightarrow f(v_2^\perp) = 1$. This can be seen as follows: $f(v_1) = 1$ implies, by the $\{0,1\}$ -coloring rules, that $f(w_i) = 0$ for all $i \in [d-2]$. Moreover, by the gadget property, we have $f(v_2) = 0$, and this imposes $f(v_2^\perp) = 1$ to satisfy the requirement that exactly one of the vertices in the maximum clique $(v_2, v_2^\perp, w_1, \dots, w_{d-2})$ is assigned value 1.

As in the original KS construction of [1], we construct a chain of $p+1$ copies $G'^{(i)}_{\text{gad}}$ ($i = 0, 1, \dots, p$) of G'_{gad} so that $p\theta = q\frac{\pi}{2}$ is an odd integral multiple of $\frac{\pi}{2}$. These copies are obtained from the realization of G'_{gad} by successive applications of a unitary \mathcal{U} , i.e., $|v_j^{(i)}\rangle = \mathcal{U}^i |v_j\rangle$ for $i = 0, 1, \dots, p$ and $j = 1, \dots, n$ and similarly for the other vectors in G'_{gad} . This unitary operator \mathcal{U} is defined as

$$\mathcal{U} = |v_2^\perp\rangle\langle v_1| - |v_2\rangle\langle v_1^\perp| + \mathbf{1}_W, \quad (6)$$

where $|v_1^\perp\rangle$ denotes the vector orthogonal to $|v_1\rangle$ in the plane $\text{span}(|v_1\rangle, |v_2\rangle)$ and where $\mathbf{1}_W$ denotes the identity on the subspace orthogonal to $\text{span}(|v_1\rangle, |v_2\rangle)$. Writing $|v_2^\perp\rangle = \alpha|v_1\rangle + \beta|v_1^\perp\rangle$ for some $\alpha, \beta \in \mathbb{C}$, we see that applying once \mathcal{U} to the faithful realization of G'_{gad} gives

$$\begin{aligned} \mathcal{U}|v_1\rangle &= |v_2^\perp\rangle, \\ \mathcal{U}|v_2^\perp\rangle &= \alpha|v_2^\perp\rangle - \beta|v_2\rangle. \end{aligned} \quad (7)$$

We have evidently $|\langle v_2^\perp | \mathcal{U} | v_2^\perp \rangle| = |\langle v_1 | v_2^\perp \rangle|$ and that

$$\arccos |\langle v_1 | \mathcal{U} | v_2^\perp \rangle| = 2 \arccos |\langle v_1 | v_2^\perp \rangle| = 2\theta. \quad (8)$$

We thus have that under successive applications of \mathcal{U} , $|v_1^{(0)}\rangle \rightarrow |v_1^{(1)}\rangle = |v_2^{\perp,(0)}\rangle$, $|v_2^{\perp,(0)}\rangle \rightarrow |v_2^{\perp,(1)}\rangle$, $|v_1^{(1)}\rangle \rightarrow |v_1^{(2)}\rangle = |v_2^{\perp,(1)}\rangle$, $|v_2^{\perp,(1)}\rangle \rightarrow |v_2^{\perp,(2)}\rangle$, and so on, with $|v_1^{(p)}\rangle \perp |v_1^{(0)}\rangle$. Furthermore, in any $\{0,1\}$ -coloring f of the graph union $\bigcup_i G'^{(i)}_{\text{gad}}$, $f(v_1^{(0)}) = 1 \Rightarrow f(v_1^{(p)}) = 1$. A similar construction of $d-1$ copies of $\bigcup_i G'^{(i)}_{\text{gad}}$ gives rise to a graph with a clique formed by the vertices $v_1^{(0)}, v_1^{(p)}$ and the $d-2$ vectors that complete the basis. The resulting graph is a Kochen-Specker graph since in any $\{0,1\}$ -coloring, if any of the vertices in this maximal clique is assigned value 1 then so are all of them, giving rise to a contradiction. We thus obtain a finite system of vectors given by the union of the vector sets in each of the graphs, that

gives rise to a proof of the Kochen-Specker theorem in dimension $\omega(G_{\text{gad}})$.

Case (ii): $\frac{\pi}{2\theta}$ is rational and is given by $\frac{p}{q}$ with q an even integer, or alternatively, $\frac{\pi}{2\theta}$ is irrational.

In this case, we construct from G'_{gad} a larger gadget \tilde{G}_{gad} with the property that the angle $\tilde{\theta}$ between the distinguished vectors obeys $\frac{\pi}{2\tilde{\theta}} = \frac{\tilde{p}}{\tilde{q}} \in \mathbb{Q}$, with \tilde{q} an odd integer. As in the previous case, we let $|v_2^\perp\rangle$ be the vector orthogonal to $|v_2\rangle$ in the plane $\text{span}(|v_1\rangle, |v_2\rangle)$, and $|v_1^\perp\rangle$ be the vector orthogonal to $|v_1\rangle$ in this plane, so that $|v_2^\perp\rangle = \alpha|v_1\rangle + \beta|v_1^\perp\rangle$, for some $\alpha, \beta \in \mathbb{C}$. We also consider a basis $\{|w_1\rangle, \dots, |w_{d-2}\rangle\}$ for the subspace orthogonal to $\text{span}(|v_1\rangle, |v_2\rangle)$ and denote G'_{gad} as the orthogonality graph of the set of vectors $\{|v_i\rangle\}_{i=1}^n \cup \{|v_2^\perp\rangle, |w_1\rangle, \dots, |w_{d-2}\rangle\}$.

Let \mathcal{U} denote a unitary operator transforming $|v_1\rangle$ to $|v_2^\perp\rangle$, i.e., \mathcal{U} is of the form

$$\begin{aligned} \mathcal{U} &= |v_2^\perp\rangle\langle v_1| - |v_2\rangle\langle v_1^\perp| + \\ &+ |w_1'\rangle\langle w_1| + \dots + |w_{d-2}'\rangle\langle w_{d-2}| \end{aligned} \quad (9)$$

with $|v_2'\rangle, |w_1'\rangle, \dots, |w_{d-2}'\rangle$ orthogonal to $|v_2^\perp\rangle$ and orthogonal to each other. Applying \mathcal{U} to the orthogonal representation of the gadget gives that

$$\begin{aligned} \mathcal{U}|v_1\rangle &= |v_2^\perp\rangle, \\ \mathcal{U}|v_2^\perp\rangle &= \alpha|v_2^\perp\rangle - \beta|v_2\rangle \end{aligned} \quad (10)$$

Let $\tilde{\theta} = \arccos |\langle v_1 | \mathcal{U} | v_2^\perp \rangle|$. We choose $|v_2'\rangle$ and thereby \mathcal{U} such that $\frac{\pi}{2\tilde{\theta}} = \frac{\tilde{p}}{\tilde{q}} \in \mathbb{Q}$ with \tilde{q} an odd integer. Now construct G'_{gad} as the orthogonality graph of the set of vectors

$$\begin{aligned} &\{|v_i\rangle\}_{i=1}^n \cup \{|v_2^\perp\rangle, |w_1\rangle, \dots, |w_{d-2}\rangle\} \cup \\ &\{\mathcal{U}|v_i\rangle\}_{i=2}^n \cup \{\mathcal{U}|v_2^\perp\rangle, |w_1'\rangle, \dots, |w_{d-2}'\rangle\}. \end{aligned} \quad (11)$$

We have thus concatenated two gadgets to form the new gadget G'_{gad} with the property that if $f(|v_1\rangle) = 1$ then also $f(|v_2^\perp\rangle) = 1$ and consequently also $f(\mathcal{U}|v_2^\perp\rangle) = 1$. We are now in the same position as in the previous case i.e., we may construct a chain of $\tilde{p}+1$ copies $G'^{(i)}_{\text{gad}}$ of G'_{gad} and follow the steps as in the previous case to construct the entire KS set in dimension $\omega(G_{\text{gad}})$.

In both cases, we thus obtain a construction of a Kochen-Specker set in dimension $\omega(G_{\text{gad}})$, completing the proof. \square

We remark that the above Theorem does not guarantee that the 01-gadgets appear as *induced* subgraphs in KS graphs; this is the case only when every vertex in the $\{0,1\}$ -edge-critical subgraph of the KS graph does not belong to three or more maximum cliques (cases (i), (ii) and (iiia) in the proof). As such, in the case (iiib) where every vertex in the $\{0,1\}$ -edge-critical subgraph of the KS graph belongs to at least three maximum cliques, the subgraphs may not correspond to vector subsets of the original KS vector set. We leave it as an interesting open question whether 01-gadgets always appear as vector subsets of

the KS vector sets in this case as well. We also note that constructions similar to that given in the proof of the second part of Theorem 1 have appeared in [33].

4 Other 01-gadgets and KS sets constructions

In this section, we make some interesting observations about 01-gadgets and provide new constructions of 01-gadgets that will be used in the next sections.

Lemma 1. *For any $d \geq 3$, there exists a 01-gadget in dimension d consisting of $5 + d$ vertices.*

Proof. For $d = 3$, a 8-vertex 01-gadget is simply given by the Clifton gadget G_{Clif} . In higher dimensions, a new 01-gadget G'_{Clif} can be obtained by adding $d - 3$ vertices to G_{Clif} with edges joining the additional vertices to each other and to each of the 8 vertices in G_{Clif} . Clearly, a faithful representation of G'_{Clif} can be obtained by supplementing the 3-dimensional representation of G_{Clif} with $d - 3$ mutually orthogonal vectors in the complementary subspace. The construction preserves the property that a $\{0, 1\}$ -coloring of G'_{Clif} exists and that the two distinguished vertices v_1, v_2 of G_{Clif} , now viewed as vertices of G'_{Clif} , cannot both be assigned the value 1 in any $\{0, 1\}$ coloring. \square

The 8-vertex Clifton gadget G_{Clif} was shown to be the minimal 01-gadget in dimension 3 [15]. This result was obtained by an exhaustive search over all non-isomorphic square-free graphs of up to 7 vertices. It is an open question to prove if the simple construction in Lemma 1 gives the minimal 01-gadgets in dimension $d > 3$ or whether even smaller gadgets exist in these higher dimensions.

In the Clifton gadget G_{Clif} the overlap between the two distinguished vertices is $|\langle v_1 | v_2 \rangle| = 1/3$. The following Lemma shows that one can reduce this overlap at the expense of increasing the dimension by one.

Lemma 2. *Let G be a 01-gadget in dimension d with distinguished vectors $|u_1\rangle, |u_2\rangle$. Then there exists a 01-gadget G' in dimension $d + 1$ with distinguished vertices $|v_1\rangle, |v_2\rangle$ for any choice of the overlap $0 < |\langle v_1 | v_2 \rangle| \leq |\langle u_1 | u_2 \rangle|$.*

Proof. Let $\{|u_i\rangle\}_{i=1}^n \subset \mathbb{C}^d$ be the set of n vectors forming the gadget G . We define G' as the set of $n + 1$ vectors $\{|v_i\rangle\}_{i=0}^n$ in \mathbb{C}^{d+1} defined as follows. For given $|u_i\rangle \in \mathbb{C}^d$, let $|\tilde{u}_i\rangle \in \mathbb{C}^{d+1}$ be the vector obtained by padding a 0 to the end of $|u_i\rangle$. Define the vectors $|v_i\rangle$ as

$$|v_i\rangle := \begin{cases} (0, \dots, 0, 1)^T, & \text{for } i = 0 \\ \mathcal{N}(|\tilde{u}_1\rangle + x(0, \dots, 0, 1)^T), & \text{for } i = 1 \\ |\tilde{u}_i\rangle & \text{for } i = 2, \dots, n \end{cases}$$

with a free parameter $x \in \mathbb{R}$ and corresponding normalization factor \mathcal{N} . Now, notice that the orthogonality relations between the set of vectors $|v_1\rangle, \dots, |v_n\rangle$ is the same as the orthogonality relations between the set of vectors $|u_1\rangle, \dots, |u_n\rangle$. The only additional orthogonality relations in G' involve $|v_0\rangle$, which is orthogonal to all other vectors but $|v_1\rangle$. By this property, it follows that if $f(|v_0\rangle) = 0$ in a coloring of G' , then the coloring of the remaining vectors $|v_1\rangle, \dots, |v_n\rangle$ is constrained exactly as for $|u_1\rangle, \dots, |u_n\rangle$ in G . In particular, we cannot have simultaneously $f(|v_1\rangle) = f(|v_2\rangle) = 1$. Now simply observe that if $f(|v_2\rangle) = 1$, we must have necessarily have $f(|v_0\rangle) = 0$ since $|v_0\rangle \perp |v_2\rangle$ and thus $|v_1\rangle$ cannot also satisfy $f(|v_1\rangle) = 1$. In other words, G' is a 01-gadget with $|v_1\rangle, |v_2\rangle$ playing the role of the distinguished vertices. Finally, we see that by varying the free parameter $x \in \mathbb{R}$, we get any overlap $0 < |\langle v_1 | v_2 \rangle| \leq |\langle u_1 | u_2 \rangle|$ between the distinguished vertices. \square

We now show the following.

Theorem 2. *Let $|v_1\rangle$ and $|v_2\rangle$ be any two distinct non-orthogonal vectors in \mathbb{C}^d with $d \geq 3$. Then there exists a 01-gadget in dimension d with $|v_1\rangle$ and $|v_2\rangle$ being the two distinguished vertices.*

While the existence of such a construction can be anticipated from the Kochen-Specker construction from Theorem 1, we give a construction with much fewer vectors based on the 43-vertex graph of Fig. 3.

Proof. The construction is based on the 43-vertex graph G of Fig. 3. We first show the construction for \mathbb{C}^3 , and then straightforwardly extend it to \mathbb{C}^d for $d > 3$. Suppose thus that we are given $|v_1\rangle, |v_2\rangle \in \mathbb{C}^3$. We consider two cases: (i) $0 < |\langle v_1 | v_2 \rangle| \leq \frac{1}{\sqrt{2}}$ and (ii) $\frac{1}{\sqrt{2}} < |\langle v_1 | v_2 \rangle| \leq 1$.

Case (i): $0 < |\langle v_1 | v_2 \rangle| \leq \frac{1}{\sqrt{2}}$. Suppose without loss of generality that $|v_1\rangle = (1, 0, 0)^T$ and $|v_2\rangle = \frac{1}{\sqrt{1+x^2}}(x, 1, 0)^T$ with $0 < x \leq 1$. In this case, the induced subgraph G_{ind} of G consisting of the vertex set $V(G_{\text{ind}}) = \{1, \dots, 22\}$ and $E(G_{\text{ind}}) = \{(u_i, u_j) : 1 \leq i, j \leq 22, (u_i, u_j) \in E(G)\}$ will suffice to construct the gadget with u_1 and u_{22} the two distinguished vertices, corresponding to $|v_1\rangle$ and $|v_2\rangle$. First, it is easily verified from the graph that in any $\{0, 1\}$ -coloring f , $f(u_1)$ and $f(u_{22})$ cannot both be assigned the value 1. It thus only remains to provide an orthogonal representation of the graph G_{ind} . Such a representation is given by the following set of (non-normalized) vectors:

$$\begin{aligned} |u_1\rangle &= (1, 0, 0)^T; & |u_2\rangle &= (0, 1, -1)^T; & |u_3\rangle &= (0, 1, 0)^T; \\ |u_4\rangle &= (0, y, 1)^T; & |u_5\rangle &= (2x, 1, 1)^T; & |u_6\rangle &= (-1, 0, 2x)^T; \\ |u_7\rangle &= (-2x, 0, -1)^T; & |u_8\rangle &= (x, 1, -2x^2)^T; \\ |u_9\rangle &= (2x^3, 2x^2, 1 + x^2)^T; \\ |u_{10}\rangle &= (-(1 + x^2), 0, 2x^3)^T; \end{aligned}$$

$$\begin{aligned}
|u_{11}\rangle &= (2x^3, 0, 1+x^2)^T; \\
|u_{12}\rangle &= (x(1+x^2), 1+x^2, -2x^4)^T; \\
|u_{13}\rangle &= (2x^5, 2x^4, (1+x^2)^2)^T; \\
|u_{14}\rangle &= (-(1+x^2)^2, 0, 2x^5)^T; \\
|u_{15}\rangle &= (2x^5, 0, (1+x^2)^2)^T; \\
|u_{16}\rangle &= (x(1+x^2)^2, (1+x^2)^2, -2x^6)^T; \\
|u_{17}\rangle &= (2x^7, 2x^6, (1+x^2)^3)^T; \\
|u_{18}\rangle &= (-x(1+y^2), -1, y)^T; \\
|u_{19}\rangle &= (1, -x, -x)^T; \quad |u_{20}\rangle = (1, -x, 0)^T; \\
|u_{21}\rangle &= (1, -x, xy)^T; \quad |u_{22}\rangle = (x, 1, 0)^T;
\end{aligned}$$

with

$$y = \frac{(1+x^2)^3 + \sqrt{(1+x^2)^6 - 16x^{14}(1+x^2)}}{4x^8}. \quad (13)$$

It is easily verified that this set of vectors satisfy all the orthogonality relations encoded by the induced subgraph G_{ind} we are considering.

Case (ii): $\frac{1}{\sqrt{2}} < |\langle v_1 | v_2 \rangle| \leq 1$. Suppose without loss of generality that $|v_1\rangle = (1, 0, 0)^T$ and $|v_2\rangle = (1+x, 1-x, 0)^T / \sqrt{2+2x^2}$ with $0 < x \leq 1$. In this case, we consider the entire 43-vertex graph G from Fig. 3, with u_1 and u_{42} the two distinguished vertices, corresponding to $|v_1\rangle$ and $|v_2\rangle$. Again, it is easily seen that in any $\{0, 1\}$ -coloring f , $f(u_1)$ and $f(u_{42})$ cannot both be assigned the value 1. It thus only remains to provide an orthogonal representation of the graph G .

The graph G can be seen as being composed from (i) the induced subgraph G_{ind} with vertices u_1, \dots, u_{22} considered above, (ii) an isomorphic subgraph G'_{ind} with vertices $u'_1 = u_{20}, u'_2 = u_{23}, \dots, u'_{22} = u_{42}$, (iii) the vertex u_{43} connected to $u_1, u_{20}, u_{22}, u_{42}$.

The first 22 vectors u_1, \dots, u_{22} of G_{ind} are chosen as above with $x = 1$ and $y = 2 + \sqrt{2}$. The 22 vectors u'_1, \dots, u'_{22} of G'_{ind} are also obtained from the above solution, but with $0 < x \leq 1$ a free parameter, and after applying first a unitary U that maps $(1, 0, 0)$ to $(1, -1, 0) / \sqrt{2}$ and $(0, 1, 0)$ to $(1, 1, 0) / \sqrt{2}$ and leave invariant $(0, 0, 1)$. We thus have $|u_1\rangle = |v_1\rangle = (1, 0, 0)^T$ and $|u_{42}\rangle = |v_2\rangle = (1+x, 1-x, 0)^T / \sqrt{2+2x^2}$ as assumed.

By construction, the orthogonality relations of the subgraphs G_{ind} and G'_{ind} are satisfied. We also have that the vectors common to the two subgraphs are indeed identical, namely $|u_{20}\rangle = (1, -1, 0)^T$ and $|u_{22}\rangle = (1, 1, 0)^T$. Furthermore, choosing $|u_{43}\rangle = (0, 0, 1)^T$, we also have that $|u_{43}\rangle$ is orthogonal to $|u_1\rangle, |u_{20}\rangle, |u_{22}\rangle$, and $|u_{42}\rangle$ as required.

This completes the construction of the gadget for \mathbb{C}^3 . Now, one may simply consider the same set of vectors as being embedded in any \mathbb{C}^d (with additional vectors $(0, 0, 0, 1, 0, \dots, 0)^T$, $(0, 0, 0, 0, 1, 0, \dots, 0)^T$ etc.) to construct a gadget in this dimension. \square

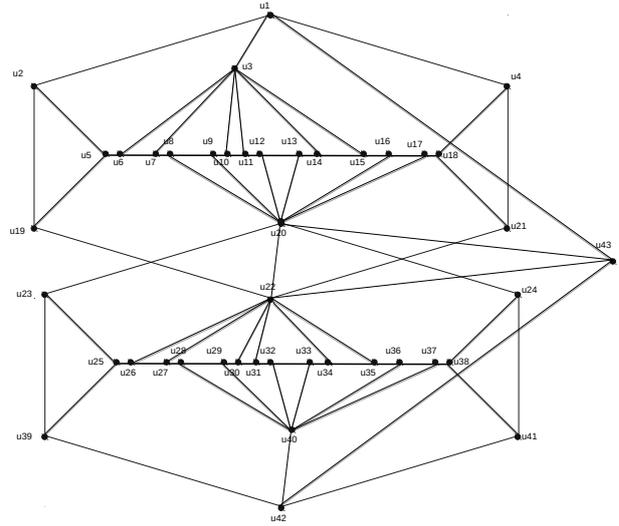


Figure 3: The 43 vertex 01-gadget used in the proof of Theorem 2.

Theorem 2 allows to construct new KS graphs than the one given in the proof of Theorem 1. Some of such constructions in dimension 3 are shown in Fig. 4. A crucial role in these is played by the repeating unit G_0 shown in Fig. 4 (a). This unit is given by a set of basis vectors $\{|u_1\rangle, |u_2\rangle, |u_3\rangle\}$ all connected via appropriate 01-gadgets to a central vector $|v_1\rangle$. In any $\{0, 1\}$ -coloring f of G_0 , one of the three basis vectors must be assigned the value 1, so that we necessarily have $f(|v_1\rangle) = 0$. In other words, G_0 is a graph in which a particular vector necessarily takes value 0 in any $\{0, 1\}$ -coloring. Note that this property is also shown by the graph in Fig. 3

Note that from G_0 , one can also construct an orthogonality graph G_1 in which a particular vector necessarily takes values 1 in any $\{0, 1\}$ -coloring. Indeed, consider two copies of G_0 with the respective central vectors $|v_1\rangle$ and $|v_2\rangle$ orthogonal to each other, so that $f(|v_1\rangle) = f(|v_2\rangle) = 0$. Then, in any $\{0, 1\}$ -coloring of the resulting graph G_1 , the third basis vector $|v_3\rangle \perp |v_1\rangle, |v_2\rangle$ necessarily obeys $f(|v_3\rangle) = 1$.

In Fig. 4 (b), a KS proof in \mathbb{C}^3 is based on the unit G_0 , repeated three times with a basis set of central vectors $|v_1\rangle, |v_2\rangle, |v_3\rangle$. By the property of G_0 in any $\{0, 1\}$ -coloring, all these three basis vectors are assigned value 0 leading to a KS contradiction. In Fig. 4 (c), the construction is based on two basis sets $\{|u_1\rangle, |u_2\rangle, |u_3\rangle\}$ and $\{|v_1\rangle, |v_2\rangle, |v_3\rangle\}$ with an appropriate 01-gadget connecting every pair $|u_i\rangle, |v_j\rangle$ for $i, j = 1, 2, 3$. So that assigning value 1 to any of the vectors in one basis, necessarily implies that all of the vectors in the other basis are assigned value 0, leading to a contradiction. Furthermore, the construction can be readily extended to derive KS graphs using any frustrated graph.

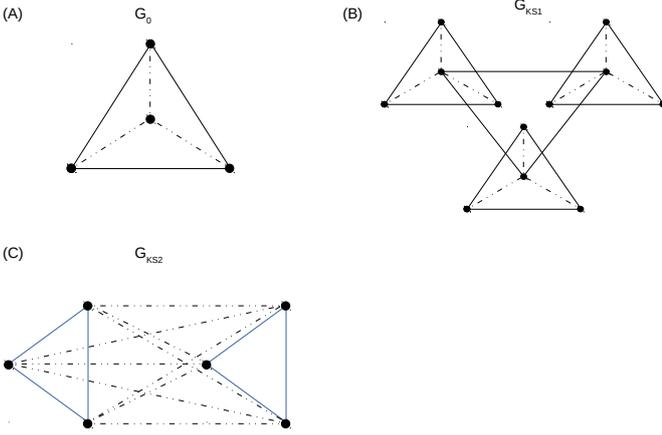


Figure 4: Graphs with the dashed edges denoting 01-gadgets. (a) In any $\{0, 1\}$ -coloring of the graph G_0 , the central vertex is necessarily assigned value 0. (b) Three copies of G_0 with the central vertices forming a basis in \mathbb{C}^3 so that the resulting graph G_{KS1} forms a Kochen-Specker proof. (c) Another proof of the KS theorem G_{KS2} is obtained by connecting every pair of vectors in two bases by a 01-gadget.

5 Statistical KS arguments based on 01-gadgets

The KS theorem can be seen as a proof that no non-contextual deterministic hidden-variable interpretation of quantum theory is possible. In a deterministic hidden-variable model, we aim to reproduce the quantum probabilities

$$\Pr_\psi(i|M) = \sum_\lambda q_\psi(\lambda) f_\lambda(i|M) \quad (14)$$

in term of hidden-variables λ , where a distribution $q_\psi(\lambda)$ over the hidden-variables is associated to each quantum state $|\psi\rangle$, and where for each λ , the model predicts with certainty that one of the outcomes i will occur for each measurement M , i.e., the hidden measurement outcome probabilities $f_\lambda(i|M)$ satisfy $f_\lambda(i|M) \in \{0, 1\}$. Furthermore, the model is non-contextual if, as in the quantum case, the probabilistic assignment to the outcome i of the (projective) measurement M , only depends on the corresponding projector V_i , independently of the wider context provided by the full description of the measurement $M = \{V_1, V_2, \dots, V_n\}$. In other words in a non-contextual deterministic hidden-variable, we aim to write for every projector V :

$$\langle \psi|V|\psi \rangle = \sum_\lambda q_\psi(\lambda) f_\lambda(V), \quad (15)$$

where $f_\lambda(V) \in \{0, 1\}$. Obviously, we should also require for consistency that $\sum_{i \in \mathcal{O}} f(V_i) \leq 1$ for any set \mathcal{O} of mutually orthogonal projectors, with equality when the projectors in \mathcal{O} sum to the identity.

No-go theorems against such models, i.e., “proofs of contextuality”, are usually obtained by considering a finite set $\mathcal{S} = \{|v_1\rangle, \dots, |v_n\rangle\} \subset \mathbb{C}^d$ of rank-one projectors V_i , represented as vectors through $V_i = |v_i\rangle\langle v_i|$. Specializing to this case, a non-contextual hidden variable model should satisfy for each $|v_i\rangle$ in \mathcal{S} and each $|\psi\rangle$ in \mathbb{C}^d ,

$$|\langle \psi|v_i\rangle|^2 = \sum_\lambda q_\psi(\lambda) f_\lambda(|v_i\rangle), \quad (16)$$

where the $f_\lambda : \mathcal{S} \rightarrow \{0, 1\}$ are $\{0, 1\}$ -colorings of \mathcal{S} .

At least three types of no-go theorems, from strongest to weakest, against such non-contextual hidden-variable models can be constructed.

The first types correspond to Kochen-Specker theorems. They establish that for certain sets \mathcal{S} , it is not possible to consistently define $\{0, 1\}$ -colorings f_λ of \mathcal{S} , even before attempting to use them to reproduce the quantum probabilities. This is what we have discussed until now.

In the second type of proofs, a $\{0, 1\}$ -coloring of \mathcal{S} is not excluded. But it can be shown that for any such coloring f_λ of \mathcal{S} , a certain inequality $\sum_i c_i f_\lambda(|v_i\rangle) \leq c_0$ must necessarily be satisfied, while in the quantum case, it happens that $\sum_i c_i |v_i\rangle\langle v_i| > c_0 \mathbb{I}$. In other words, though it is possible to find a $\{0, 1\}$ assignment $f_\lambda(|v_i\rangle)$ to each projector $|v_i\rangle\langle v_i|$ in \mathcal{S} that is compatible with the orthogonality relations among such projectors, any such assignment fails to reproduce some more complex relation of the type $\sum_i c_i |v_i\rangle\langle v_i| > c_0 \mathbb{I}$ satisfied by these projectors. This immediately implies a contradiction with eq. (16), since in the quantum case we have for any $|\psi\rangle$, $\sum_i c_i |\langle \psi|v_i\rangle|^2 > c_0$, while according to a non-contextual hidden variable model, we would have $\sum_i c_i |\langle \psi|v_i\rangle|^2 = \sum_\lambda q_\psi(\lambda) [\sum_i c_i f_\lambda(|v_i\rangle)] \leq \sum_\lambda q_\psi(\lambda) c_0 \leq c_0$. Such no-go theorems are referred to as “statistical state-independent” KS arguments and were introduced by Yu and Oh [25].

Finally, for certain sets \mathcal{S} , it is possible to find valid $\{0, 1\}$ -colorings that do not lead to any type of contradictions of the second type above. However, it is not possible to take mixtures of such colorings, as in eq. (16), to reproduce the predictions of certain quantum states $|\psi\rangle$. Such no-go theorems are referred to as “statistical state-dependent” KS arguments and were introduced by Clifton in [17].

While we have seen in the previous section how proofs of the KS theorem can be constructed using 01-gadgets, in this section we show how to use them to build statistical state-independent and state-dependent KS arguments

5.1 State-independent KS arguments

In [25], Yu and Oh introduced a set of 13 vectors in \mathbb{C}^3 that provides a state-independent proof of contextuality, despite not being a KS set. We show how

using Theorem 2, it is possible to construct other state-independent proofs of contextuality based on 01-gadgets. To do this, we make use of the following lemma.

Lemma 3. *Let $|u_i\rangle$, for $i = 1, \dots, d+1$ be the unit vectors denoting the vertices of a d -dimensional simplex embedded in \mathbb{R}^d . Then*

$$\sum_{i=1}^{d+1} |u_i\rangle\langle u_i| = \frac{d+1}{d} \mathbb{I}. \quad (17)$$

Proof. Since $|u_i\rangle$ form the vertices of the d -simplex, we have $\langle u_i|u_j\rangle = -\frac{1}{d}$ for any $i \neq j \in \{1, \dots, d+1\}$. It then follows

$$\left(\sum_{i=1}^{d+1} \langle u_i| \right) \left(\sum_{j=1}^{d+1} |u_j\rangle \right) = (d+1) + d(d+1) \left(-\frac{1}{d} \right) = 0,$$

so that

$$O := \sum_{i=1}^{d+1} |u_i\rangle\langle u_i| = - \sum_{i \neq j=1}^{d+1} |u_i\rangle\langle u_j| \quad (18)$$

This then implies that

$$O^2 = O - \frac{1}{d} \sum_{i \neq j=1}^{d+1} |u_i\rangle\langle u_j| = \frac{d+1}{d} O. \quad (19)$$

Moreover, O is invertible, since $\text{span}(\{|u_i\rangle\}_{i=1}^{d+1}) = \mathbb{R}^d$ so that we obtain $O = \frac{d+1}{d} \mathbb{I}$. \square

Now, state-independent KS arguments for \mathbb{C}^d are straightforwardly constructed as follows. For every pair of vectors $|u_i\rangle, |u_j\rangle$ of the d -simplex, consider a 01-gadget \mathcal{S}_{ij} with $|u_i\rangle, |u_j\rangle$ the distinguished vertices. Since $|u_i\rangle$ and $|u_j\rangle$ are non-orthogonals, such gadgets exists, as implied by Theorem 2. The resulting set of vectors $\mathcal{S} = \cup_{ij} \mathcal{S}_{ij}$ exhibits state-independent contextuality. Indeed, by the property of the 01-gadgets, only one of the vectors $|u_i\rangle$ for $i = 1, \dots, d+1$ can be assigned the value 1 in any $\{0, 1\}$ -coloring of \mathcal{S} . It thus follows that

$$\sum_{i=1}^{d+1} f(|u_i\rangle) \leq 1, \quad (20)$$

On the other hand, from Lemma 3, every state $|\psi\rangle$ from \mathbb{C}^d achieves the value $\sum_{i=1}^{d+1} |\langle \psi|u_i\rangle|^2 = \frac{d+1}{d} > 1$.

While we have used the $d+1$ vertices of a d -simplex in the construction above, we observe that any set $\{|u_i\rangle\}$ of vectors in \mathbb{C}^d such that $\sum_i |\langle \psi|u_i\rangle|^2 > 1$ for all $|\psi\rangle \in \mathbb{C}^d$ can be utilized in the construction, although such a set clearly needs to contain at least $d+1$ vectors.

5.2 State-dependent KS arguments

The relation between state-dependent KS arguments and 01-gadgets is even more direct than in the above construction. Actually, the first state-dependent KS argument introduced by Clifton in [17] was precisely based on the set of vectors (5) forming the Clifton gadget G_{gad} . His argument was as follows. In every non-contextual hidden-variable model attempting to replicate the quantum probabilities associated to the projectors of the Clifton gadget, we should have $|\langle \psi|u_1\rangle|^2 + |\langle \psi|u_8\rangle|^2 = \sum_{\lambda} q_{\psi}(\lambda) (f_{\lambda}(|u_1\rangle) + f_{\lambda}(|u_8\rangle)) \leq 1$, by the gadget property. However, if we take $|\psi\rangle = |u_1\rangle$, we find that according to the quantum predictions $|\langle u_1|u_1\rangle|^2 + |\langle u_1|u_8\rangle|^2 = 1 + |\langle u_1|u_8\rangle|^2 > 1$ since $|\langle u_1|u_8\rangle|^2 > 0$ as $|u_1\rangle$ and $|u_8\rangle$ are non-orthogonal. Other state-dependent proofs based on inequalities have since been developed, with the smallest involving five vectors [4]. The first state-independent statistical KS argument was presented in [20] and the proof that any KS set give can be converted in a state-independent statistical KS argument was presented in [33].

Obviously, the argument used by Clifton for the particular set of vectors he introduced, immediately carries over to any 01-gadget. Thus every 01-gadget serves as a proof of state-dependent contextuality.

Note that it was realized in [13] that a class of graphs, known as perfect graphs, define a class of graphs that cannot serve as proofs of (even state-dependent) contextuality. That is, for any orthogonal representation $\{|v_j\rangle\} \subset \mathbb{C}^d$ of a perfect graph and for any pure state $|\psi\rangle \in \mathbb{C}^d$, the outcome probabilities $|\langle \psi|v_j\rangle|^2$ admit a non-contextual hidden variable model of the form (16). Since a non-contextual hidden variable model is not possible for a 01-gadget, we deduce that no perfect graph is a 01-gadget. Perfect graphs are a well-known class of graphs which by the strong perfect graph theorem [30] can be characterized as those graphs that do not contain odd cycles and anti-cycles of length greater than three as induced subgraphs.

Finally, remark that the argument due to Clifton presented above works not only for the state $|\psi\rangle = |u_1\rangle$, but for any state $|\psi\rangle \in \mathbb{C}^3$ which obeys $|\langle \psi|u_1\rangle|^2 + |\langle \psi|u_8\rangle|^2 > 1$. More generally, we now present a 01-gadget which serves to prove state-dependent contextuality for all but a measure zero set of states in \mathbb{C}^3 .

This construction is based on the gadget G of Fig. 3 with the 43 vector orthogonal representation presented in the proof of Theorem 2. Note that if we take $x = 1$ in this representation, then the two distinguished vectors $|u_1\rangle$ and $|u_{42}\rangle$ actually coincide and are both equal to $(1, 0, 0)$ (i.e., the two distinguished vertices u_1 and u_{42} should actually be identified). Therefore in any $\{0, 1\}$ -coloring f of G , $2f(|u_1\rangle) = f(|u_1\rangle) + f(|u_{42}\rangle) \leq 1$, i.e. the vector $|v_1\rangle$ is assigned value 0. This implies that G witnesses

state-dependent contextuality of all states in \mathbb{C}^3 but for a measure zero set of states $|\psi\rangle$ that are orthogonal to $|v_1\rangle = (1, 0, 0)$.

The construction that we just described is based on 42 vectors. It is actually possible to find a slightly smaller construction based on the following 40 vectors:

$$\begin{aligned}
|u_1\rangle &= (1, -1, 0)^T; & |u_2\rangle &= (1, 1, 1)^T; \\
|u_3\rangle &= (1, 1, 0)^T; & |u_4\rangle &= (1, 1, b)^T; \\
|u_5\rangle &= (-2, 1, 1)^T; & |u_6\rangle &= (1, -1, 3)^T; \\
|u_7\rangle &= (3, -3, -2)^T; & |u_8\rangle &= (2, 0, 3)^T; \\
|u_9\rangle &= (-3, 0, 2)^T; & |u_{10}\rangle &= (-2, 2, -3)^T; \\
|u_{11}\rangle &= (3, -3, -4)^T; & |u_{12}\rangle &= (4, 0, 3)^T; \\
|u_{13}\rangle &= (-3, 0, 4)^T; & |u_{14}\rangle &= (-4, 4, -3)^T; \\
|u_{15}\rangle &= (3, -3, -8)^T; & |u_{16}\rangle &= (8, 0, 3)^T; \\
|u_{17}\rangle &= (-3, 0, 8)^T; & |u_{18}\rangle &= (-8, 4 + \sqrt{7}, -3)^T; \\
|u_{19}\rangle &= (0, 1, -1)^T; & |u_{20}\rangle &= (0, 1, 0)^T; \\
|u_{21}\rangle &= (0, -3 + 8b, -16 - 3b)^T; & |u_{22}\rangle &= (1, 0, 0)^T; \\
|u_{23}\rangle &= (1, 0, -1)^T; & |u_{24}\rangle &= (2 - \sqrt{2}, 0, 1)^T; \\
|u_{25}\rangle &= (1, -2, 1)^T; & |u_{26}\rangle &= (0, 1, 2)^T; \\
|u_{27}\rangle &= (0, 2, -1)^T; & |u_{28}\rangle &= (1, -1, -2)^T; \\
|u_{29}\rangle &= (1, -1, 1)^T; & |u_{30}\rangle &= (0, 1, 1)^T; \\
|u_{31}\rangle &= (0, 1, -1)^T; & |u_{32}\rangle &= (-1, 1, 1)^T; \\
|u_{33}\rangle &= (-1, 1, -2)^T; & |u_{34}\rangle &= (0, 2, 1)^T; \\
|u_{35}\rangle &= (0, 1, -2)^T; & |u_{36}\rangle &= (2, -2, -1)^T; \\
|u_{37}\rangle &= (1, -1, 4)^T; & |u_{38}\rangle &= (-2 - \sqrt{2}, 6 - \sqrt{2}, 2)^T; \\
|u_{39}\rangle &= |u_2\rangle; & |u_{40}\rangle &= |u_3\rangle; & |u_{41}\rangle &= (1, 1, -2 + \sqrt{2})^T; \\
|u_{42}\rangle &= |u_1\rangle; & |u_{43}\rangle &= (0, 0, 1)^T;
\end{aligned}$$

with $b = \frac{-4 + \sqrt{7}}{3}$, and where we have the following identities $|u_1\rangle = |u_{42}\rangle$, $|u_2\rangle = |u_{39}\rangle$, $|u_3\rangle = |u_{40}\rangle$. It can be verified that the graph in Fig. 3 where we identify the vertices u_1 and u_{42} , u_2 and u_{39} , u_3 and u_{40} , is the orthogonality graph of these 40 vectors. These 40 vectors thus form a 01-gadget, where as above the vector $|u_1\rangle = (1, -1, 0)$ can only be assigned the value 0, implying that it can serve as a state-dependent contextuality proof for any vector in \mathbb{C}^3 that is not orthogonal to $(1, -1, 0)$. We leave it as an open question whether this set of 40 vectors is the minimal set with this property.

6 Proofs of the extended Kochen-Specker theorem using 01-gadgets

In this section, we consider a stronger variant of the KS theorem due to Pitowsky [22] and Hrushovski and Pitowsky [23]. While the KS theorem is concerned with $\{0, 1\}$ -colorings where all projectors (or vectors) in a given set S must be assigned a value in $\{0, 1\}$, we consider here more general assignments where any real

value in $[0, 1]$ is allowed to the members of S . Specifically, given a set of vectors $\mathcal{S} = \{|v_1\rangle, \dots, |v_n\rangle\} \subset \mathbb{C}^d$, we say that $f : \mathcal{S} \rightarrow [0, 1]$ is a $[0, 1]$ -assignment if f satisfies the same rules (3) as it does for $\{0, 1\}$ -colorings. Both $\{0, 1\}$ -colorings and $[0, 1]$ -assignments can be interpreted as assigning a probability to the projectors corresponding to each of the elements of S . But while the assignment is constrained to be deterministic in the case of $\{0, 1\}$ -colorings since these probabilities can only take the values 0 or 1, the probabilistic assignment may be completely general (hence non-deterministic) for $[0, 1]$ -assignments. In particular, for any given quantum state $|\psi\rangle$, the Born rule $f(|v_i\rangle) = |\langle\psi|v_i\rangle|^2$ defines a valid $[0, 1]$ -assignment.

Hrushovski and Pitowsky [23], following earlier work by Pitowsky in [22], proved the following theorem, which they call the “logical indeterminacy principle”.

Theorem 3 ([23]). *Let $|v_1\rangle$ and $|v_2\rangle$ be two non-orthogonal vectors in \mathbb{C}^d with $d \geq 3$. Then there is a finite set of vectors $S \subset \mathbb{C}^d$ with $|v_1\rangle, |v_2\rangle \in S$ such that for any $[0, 1]$ -assignment, it holds that $f(|v_1\rangle), f(|v_2\rangle) \in \{0, 1\}$ if and only if $f(|v_1\rangle) = f(|v_2\rangle) = 0$.*

Thus for any two non-orthogonal vectors $|v_1\rangle$ and $|v_2\rangle$, at least one of the probabilities associated to the vectors $|v_1\rangle$ or $|v_2\rangle$ must be strictly between zero and one, unless they are both equal to zero. A corollary of this result, observed in [3, 9, 8] is that if $f(|v_1\rangle) = 1$ (this should, for instance, necessarily be the case if we attempt to reproduce the quantum probabilities for measurements performed on the state $|\psi\rangle = |v_1\rangle$), then $f(|v_2\rangle) \neq 0, 1$, showing that one can localise the “value-indefiniteness” of quantum observables that the KS theorem implies. Theorem 3 therefore provides a stronger variant of the KS theorem, and we will refer to it as the *extended KS theorem*.

The proof of Theorem 3 given in [23] was obtained as a corollary of Gleason’s theorem [24]. A more explicit constructive proof was given by Abbott, Calude and Svozil [3, 9], where they also noted that significantly none of the known KS sets serves to prove Theorem 3. Note that an earlier proof of the extended KS theorem was also given in [22]. All these existing proofs of the extended KS theorem involve complicated constructions with no systematic procedure for obtaining the requisite sets of vectors. In this subsection, we will provide a simple systematic method for obtaining in a constructive way these extended KS sets.

In order to prove the extended KS theorem, we need gadgets of a special kind, which are defined as usual 01-gadgets apart from the fact that the condition that the two distinguished vertices cannot both be assigned the value 1 in any $\{0, 1\}$ -colorings should also hold for any $[0, 1]$ -assignments. That is, we simply replace ‘ $\{0, 1\}$ -coloring’ by ‘ $[0, 1]$ -assignment’ and

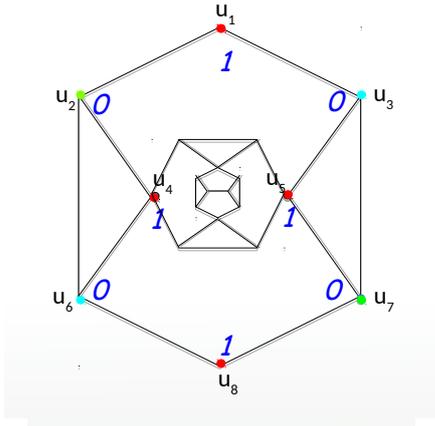


Figure 5: An iterative construction of an extended 01-gadget for which the two distinguished vertices u_1 and u_8 are such that in the limit of large number of iterations k , $|\langle u_1^{(k)} | u_8^{(k)} \rangle| \in [0, 1]$.

$f(|v_1\rangle) + f(|v_2\rangle) \leq 1$ by $f(|v_1\rangle) + f(|v_2\rangle) < 2$ in Definition 1, and similarly for Definition 2. We call such new gadgets ‘extended 01-gadgets’. It is easily verified that the Clifton gadget in Fig. 1 and the 16-vertex gadget in Fig. 2 obey this additional restriction.

Our first aim will be to construct such extended 01-gadgets for any two given non-orthogonal vectors $|v_1\rangle, |v_2\rangle \in \mathbb{C}^d$ for $d \geq 3$. This is the content of the following Theorem, which generalizes Theorem 2.

Theorem 4. *Let $|v_1\rangle$ and $|v_2\rangle$ be any two distinct non-orthogonal vectors in \mathbb{C}^d with $d \geq 3$. Then there exists an extended 01-gadget in dimension d with $|v_1\rangle$ and $|v_2\rangle$ being the two distinguished vertices.*

Proof. We begin with the construction for $d = 3$ and generalize it to higher dimensions naturally. The construction is an iterative procedure based on the Clifton gadget G_{Clif} given in Fig. 1.

Firstly, as stated previously, it is readily seen that G_{Clif} is actually an extended 01-gadget with u_1, u_8 the two distinguished vertices, i.e., any $[0, 1]$ -assignment $f : V(G_{\text{Clif}}) \rightarrow [0, 1]$ cannot be such that $f(u_1) = f(u_8) = 1$. Further, it is known that the \mathbb{R}^3 realization of G_{Clif} given by (5) achieves the (minimal possible) separation of $\theta_1 = \arccos |\langle u_1 | u_8 \rangle| = \arccos 1/3$ between the two end vertices [19].

We now describe a nesting procedure that at each

step decreases the angle between the vectors corresponding to the two outer vertices. The procedure works as follows. Replace the edge (u_4, u_5) in G_{Clif} by G'_{Clif} , a copy of G_{Clif} where we identify $u'_1 = u_4$ and $u'_8 = u_5$. The new graph thus obtained has 14 vertices and 21 edges. The operation has the property that in any $[0, 1]$ -assignment f , an assignment of value 1 to the two outer vertices of the new graph (i.e. u_1, u_8) leads to a similar assignment to the two outer vertices of the inner copy of G_{Clif} (i.e. u'_1, u'_8) thereby giving rise to a contradiction. In other words, the newly constructed graph is once again an extended 01-gadget. This procedure can be repeated an arbitrary number of times, as illustrated in Fig. 5, leading to an extended 01-gadget formed from k nested Clifford graphs $G^1_{\text{Clif}}, G^2_{\text{Clif}}, G^3_{\text{Clif}}, \dots, G^k_{\text{Clif}}$ where G^1_{Clif} corresponds to the most inner graph and G^k_{Clif} to the most outer graph. We now show that the total graph at the k -th iteration is an orthogonality graph where the overlap $|\langle u_1^{(k)} | u_8^{(k)} \rangle|$ between the two outer vertices u_1^k, u_8^k can be chosen to take any value in $[0, \frac{k}{k+2}]$, thus spanning any possible value in $[0, 1]$ for k sufficiently large. Setting $|v_1\rangle = |u_1^{(k)}\rangle$ and $|v_2\rangle = |u_8^{(k)}\rangle$ with k depending on the overlap of the given vectors $|\langle v_1 | v_2 \rangle|$, then gives the required gadget and proves the Theorem.

Suppose that at the k -th step of the iteration, the vectors representing the two outer vertices of the ‘inner’ gadget from the $k - 1$ -th step are

$$\begin{aligned} |u_4^{(k)}\rangle &= |u_1^{(k-1)}\rangle = (1, 0, 0), \\ |u_5^{(k)}\rangle &= |u_8^{(k-1)}\rangle = \frac{1}{\sqrt{1+x_k^2}}(x_k, 1, 0), \end{aligned} \quad (21)$$

without loss of generality, so that the overlap between these vectors is $|\langle u_4^{(k)} | u_5^{(k)} \rangle| = \frac{x_k}{\sqrt{1+x_k^2}}$, where for simplicity of the construction we take $x_k \in \mathbb{R}_0^+$. The remaining vectors then in general have the following (non-normalized) orthogonal representation in \mathbb{R}^3

$$\begin{aligned} |u_8^{(k)}\rangle &= (a_k, b_k, c_k), \quad |u_6^{(k)}\rangle = (0, -c_k, b_k), \\ |u_7^{(k)}\rangle &= (c_k, -c_k x_k, -a_k + b_k x_k), \quad |u_2^{(k)}\rangle = (0, b_k, c_k), \\ |u_3^{(k)}\rangle &= (-a_k + b_k x_k, a_k x_k - b_k x_k^2, -c_k - c_k x_k^2), \\ |u_1\rangle &= (-b_k c_k - a_k c_k x_k, -a_k c_k + b_k c_k x_k, a_k b_k - b_k^2 x_k), \end{aligned} \quad (22)$$

with $a_k, b_k, c_k \in \mathbb{R}$. This gives an overlap of

$$|\langle u_1^{(k)} | u_8^{(k)} \rangle| = \frac{|-a_k c_k (b_k + a_k x_k)|}{\sqrt{(a_k^2 + b_k^2 + c_k^2)(c_k^2 (b_k + a_k x_k)^2 + b_k^2 (a_k - b_k x_k)^2 + (a_k c_k - b_k c_k x_k)^2)}}. \quad (23)$$

A direct optimization of this expression with respect to the parameters a_k, b_k, c_k gives the choice $b_k = 1, c_k = 1, a_k = x_k + \sqrt{1+x_k^2}$. So that the overlap

between the two outer vertices at the k -th step of the

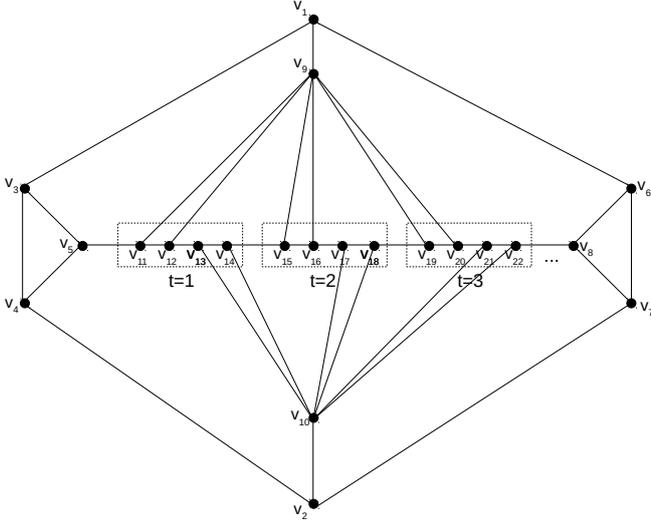


Figure 6: An alternative construction of an extended 01-gadget for which the two distinguished vertices v_1 and v_2 are such that in the limit of large number t of the repeating unit of four vectors, $|\langle v_1 | v_2 \rangle|$ can take any value in $[0, 1]$.

iteration is given by

$$|\langle u_1^{(k)} | u_8^{(k)} \rangle| = \frac{1}{3 + 4x_k(x_k - \sqrt{1+x_k^2})} =: \frac{x_{k+1}}{\sqrt{1+x_{k+1}^2}} \quad (24)$$

With the initial overlap for $k = 1$ of $1/3$ and corresponding initial x values of $x_1 = 0$ and $x_2 = \frac{1}{2\sqrt{2}}$, we can now evaluate the expression for the overlap for any $k > 1$. We find that the overlap at the k -th step is $\frac{k}{k+2}$. This is readily seen by an inductive argument. The base claim is clear, suppose that at the k -th step the overlap is given by $\frac{x_{k+1}}{\sqrt{1+x_{k+1}^2}} = \frac{k}{k+2}$, i.e., $x_{k+1} = \frac{k}{2\sqrt{k+1}}$. Substituting in Eq.24, we obtain $\frac{x_{k+2}}{\sqrt{1+x_{k+2}^2}} = \frac{k+1}{k+3} = \frac{(k+1)}{(k+1)+2}$. Moreover, we see that choosing $b_k = 1, c_k = 1$, the overlap expression (23) is a continuous function of a_k for any fixed x_k with the minimum value of 0 achieved at $a_k = 0$. Thus, every intermediate overlap in $[0, \frac{k}{k+2}]$ between the two outer vectors is also achievable by appropriate choice of a_k for the fixed value of x_k, b_k, c_k . This completes the construction of the gadget for \mathbb{C}^3 (possibly by taking its faithful version in the graph representation).

Now, one may simply consider the same set of vectors as being embedded in any \mathbb{C}^d (with additional vectors $(0, 0, 0, 1, 0, \dots, 0)^T$, $(0, 0, 0, 0, 1, 0, \dots, 0)^T$ etc.) to construct a gadget in this dimension. \square

In fact, the construction above is not unique. We give an alternative set of vectors that also serves to prove Theorem 4. The construction is shown in Fig. 6. Suppose we are given two distinct non-orthogonal vectors $|v_1\rangle = (1, 0, 0)^T$, $|v_2\rangle = (x, \sqrt{1-x^2}, 0)^T$, with $0 < x < 1$. We begin by adding the following set of

vectors with a parameter $y \in \mathbb{R}$:

$$\begin{aligned} |v_3\rangle &= (0, x, -\sqrt{1-x^2})^T; \\ |v_4\rangle &= (-(1-x^2), x\sqrt{1-x^2}, x^2)^T; \\ |v_5\rangle &= (x, (1-x^2)\sqrt{1-x^2}, x(1-x^2))^T; \\ |v_6\rangle &= (0, y, \sqrt{1-y^2})^T; \\ |v_7\rangle &= (-\sqrt{(1-x^2)(1-y^2)}, x\sqrt{1-y^2}, xy)^T; \\ |v_8\rangle &= (x, (1-y^2)\sqrt{1-x^2}, y\sqrt{(1-x^2)(1-y^2)})^T; \\ |v_9\rangle &= (0, 1, 0)^T; \quad |v_{10}\rangle = (-\sqrt{1-x^2}, x, 0)^T. \end{aligned} \quad (25)$$

The remaining vectors are obtained using a repeating unit consisting of four vectors:

$$\begin{aligned} |v_{7+4t}\rangle &= (-(1-x^2), 0, x^{2(t-1)})^T; \\ |v_{8+4t}\rangle &= (x^{2(t-1)}, 0, 1-x^2)^T; \\ |v_{9+4t}\rangle &= (-x(1-x^2), -(1-x^2)\sqrt{1-x^2}, x^{2t-1})^T; \\ |v_{10+4t}\rangle &= (x^{2t}, x^{2t-1}\sqrt{1-x^2}, 1-x^2)^T; \end{aligned} \quad (26)$$

repeated t times for an integer $t \geq 1$ depending on x . Choosing the parameter y as

$$y = \sqrt{\frac{(1-x^2)^2 + 2x^{4t-2} - \sqrt{(1-x^2)((1-x^2)^3 - 4x^{4t}})}{2(1-x^2)(1-x^2 + x^{4t-2})}},$$

we find that $y \in \mathbb{R}$, for t satisfying $(1-x^2)^3 \geq 4x^{4t}$. We see that as t increases this inequality can be satisfied for larger values of x , and for any $0 < x < 1$ as $t \rightarrow \infty$. From the orthogonality graph of this set of vectors S shown in Fig. 6, it is clear that there cannot be any assignment $f : S \rightarrow [0, 1]$ such that $f(|v_1\rangle) = f(|v_2\rangle) = 1$, giving an extended 01-gadget.

While the construction in Theorem 4 and that in the previous paragraph work for any two distinct vectors, given two such vectors it is of great interest to find the minimal extended 01-gadget with these vectors as the distinguished vertices. While this question is the foundational analog for extended KS systems of the question of finding minimal KS sets, it is also of practical interest in obtaining Hardy paradoxes with optimal values of the non-zero probability, and extracting randomness from the gadgets [29].

We now show how the extended 01-gadgets can be used to construct proofs of the extended KS Theorem 3.

Proof. (Theorem 3) We present the construction for $d = 3$, the proof for higher dimensions will follow in an analogous fashion. The idea is encapsulated by Fig. 7. Suppose we are given two distinct non-orthogonal vectors $|v_1\rangle$ and $|v_2\rangle$ in \mathbb{C}^d . We begin by constructing an appropriate extended 01-gadget G_{v_1, v_2} , depending on $|\langle v_1 | v_2 \rangle|$, with the corresponding v_1, v_2 being the distinguished vertices.

Let $|v_3\rangle = |v_1\rangle \times |v_2\rangle$ denote the vector orthogonal to the plane $\text{span}(|v_1\rangle, |v_2\rangle)$ spanned by $|v_1\rangle$ and $|v_2\rangle$, where \times denotes the cross product of the vectors. Let $|v_4\rangle$ be the vector in the

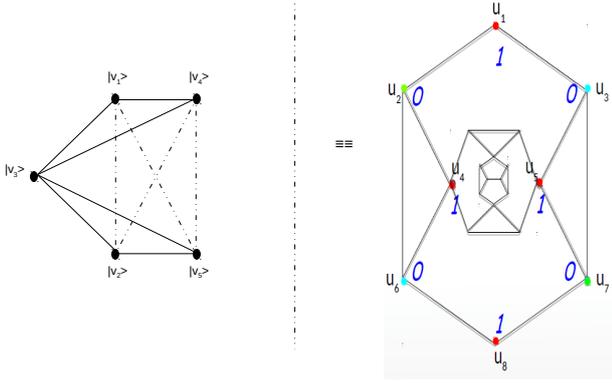


Figure 7: A constructive proof of the extended Kochen-Specker theorem 3 using the extended 01-gadgets. Given vectors $|v_1\rangle, |v_2\rangle \in \mathbb{C}^d$, we obtain vector $|v_3\rangle \perp \text{span}(|v_1\rangle, |v_2\rangle)$ and two other vectors $|v_4\rangle, |v_5\rangle$ in the plane $\text{span}(|v_1\rangle, |v_2\rangle)$ with the orthogonality relations indicated in the left figure. Dashed edges between two vertices indicate an extended 01-gadget from Theorem 4 with the two vertices being distinguished.

plane $\text{span}(|v_1\rangle, |v_2\rangle)$ orthogonal to $|v_1\rangle$, and $|v_5\rangle$ denote the vector in this plane orthogonal to $|v_2\rangle$, so that $\{|v_1\rangle, |v_3\rangle, |v_4\rangle\}$, $\{|v_2\rangle, |v_3\rangle, |v_5\rangle\}$ form orthogonal bases in \mathbb{C}^3 . We construct appropriate extended 01-gadgets G_{v_1, v_5} and G_{v_2, v_4} depending on $|\langle v_1 | v_5 \rangle|$ and $|\langle v_2 | v_4 \rangle|$. In G_{v_1, v_5} the vertices v_1, v_5 corresponding to the vectors $|v_1\rangle, |v_5\rangle$ play the role of the distinguished vertices and similarly in G_{v_2, v_4} . Let G_{Pit} denote the orthogonality graph of the entire set of vectors $G_{v_1, v_2} \cup G_{v_1, v_5} \cup G_{v_2, v_4} \cup |v_3\rangle$.

We have that in any assignment $f : V(G_{\text{Pit}}) \rightarrow [0, 1]$ for which $f(v_1), f(v_2) \in \{0, 1\}$, if $f(v_1) = 1, f(v_2) = 1$, then we obtain a contradiction by the property of the extended 01-gadget G_{v_1, v_2} . On the other hand, if $f(v_1) = 1, f(v_2) = 0$, then since $|v_1\rangle \perp |v_3\rangle$ we have $f(v_3) = 0$, and by the property of the extended 01-gadget G_{v_1, v_5} we have $f(v_5) = 0$. This gives a contradiction since v_2, v_3, v_5 form a maximum clique. Similarly, if $f(v_1) = 0, f(v_2) = 1$, then since $|v_2\rangle \perp |v_3\rangle$ we have $f(v_3) = 0$, and by the property of the extended 01-gadget G_{v_2, v_4} we have $f(v_4) = 0$. This also gives a contradiction since v_1, v_3, v_4 form a maximum clique. Therefore, we have any assignment $f : V(G_{\text{Pit}}) \rightarrow [0, 1]$ which obeys $f(v_1), f(v_2) \in \{0, 1\}$ also must necessarily obey $f(v_1) = f(v_2) = 0$. This completes the proof. \square

6.1 Discussion

Intuitively, with respect to any $\{0, 1\}$ coloring, a 01-gadget behaves like a "virtual edge" between its two special vertices, with this edge also obeying the rule

that at most one of its incident vertices may be assigned the color 1. Moreover, in Theorem 2 we have shown that 01-gadgets may be constructed with any two non-orthogonal vectors as the special vertices. Starting from a given set of vectors, this allows us to connect any two non-orthogonal vectors by an appropriate 01-gadget, which imposes additional constraints on the $\{0, 1\}$ -colorings of the resulting set of vectors. By appropriately adding such virtual edges, we are eventually able to obtain a set of vectors that gives a Kochen-Specker contradiction. Moreover, it turns out that the statistical proofs of the Kochen-Specker theorem can also be interpreted in the same manner. For instance, the famous Yu-Oh graph of [25] can be interpreted as six 01-gadgets connecting the vectors $(1, 1, 1)^T, (1, 1, -1)^T, (1, -1, 1)^T$ and $(-1, 1, 1)^T$. These four vectors thus form a "virtual clique", with the property that in any $\{0, 1\}$ -coloring of the Yu-Oh set, the sum of the values attributed to these four vectors cannot exceed one. On the other hand, any quantum state has overlap with these four vectors summing to $4/3$ providing a statistical contradiction. Similar considerations also apply to the extended Kochen-Specker theorem of Pitowsky by means of extended 01-gadgets.

7 Computational complexity of $\{0, 1\}$ -colorings

Clearly, complete graphs of size $d+1$ cannot be faithfully realized in \mathbb{C}^d , but there also exist certain other graphs that cannot be faithfully realized in \mathbb{C}^d . The well-known example is the four-cycle (square) graph in \mathbb{C}^3 , this can be seen by the following simple argument. Suppose a pair of vertices in opposite corners of the square is assigned without loss of generality the vectors $|0\rangle$ and $\alpha|0\rangle + \beta|1\rangle$, with $\alpha, \beta \in \mathbb{C}$. Since these vectors span a plane and the remaining pair of vertices are both required to be orthogonal to this plane, these latter vectors are both equal up to a phase to $|2\rangle$, contradicting the requirement of faithfulness. There exist analogous graphs that are not faithfully realizable in higher dimensions, some of which are shown in Fig. 8. In searching for Kochen-Specker vector systems in \mathbb{C}^d , it is therefore crucial to reduce the size of the search by restricting to non-isomorphic graphs which do not contain these forbidden graphs as subgraphs. Indeed, searching over non-isomorphic square-free graphs lead to the proof that the smallest Kochen-Specker vector system in \mathbb{C}^3 is of size at least 18 [15].

Let us denote the set of forbidden graphs in \mathbb{C}^d as $\{G_{\text{fbd}}\}$. We show, following the proof by Arends et al. [15, 16] for the square-free case, that the problem of checking $\{0, 1\}$ -colorability of $\{G_{\text{fbd}}\}$ -free graphs is NP-complete. Here, by a $\{G_{\text{fbd}}\}$ -free graph we mean a graph that does not contain any of the forbidden graphs as subgraphs.

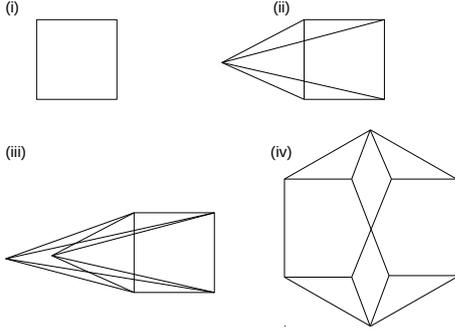


Figure 8: Examples of forbidden subgraphs in dimensions 3, 4 and 5. Graph (i) is the square graph which is not faithfully realizable in \mathbb{C}^3 as explained in the text. Graph (iv) is the graph from [14] which was verified to be not faithfully realizable in dimension three despite being square-free. Graph (ii) is not faithfully realizable in \mathbb{C}^4 , which can be seen as arising from the fact that the induced square subgraph is not faithfully realizable in \mathbb{C}^3 and the additional vertex being adjacent to all vertices of the square, the vector corresponding to this vertex occupies an orthogonal subspace to that spanned by the square. Graph (iii) is similarly not realizable in \mathbb{C}^5 this time owing to the presence of two vertices (which themselves cannot be represented by identical vectors) that are adjacent to all the vertices of the square. It is clear that the construction can be extended to higher dimensions.

Theorem 5 (see also [15]). *Checking $\{0,1\}$ -colorability of $\{G_{\text{fbd}}\}$ -free graphs is NP-complete.*

The proof is based on a reduction to the well-known graph coloring problem that uses 01-gadgets in a crucial manner. Let us first recall the usual notion of coloring of a graph used in the proof. A proper coloring c of a graph G is an assignment of one among n colors to each of the vertices of the graph $c : V(G) \rightarrow [n]$ ($[n] := \{1, \dots, n\}$) such that no pair of adjacent vertices are assigned the same color. If such a coloring exists, we say that G is n -colorable.

Proof. The proof generalizes and simplifies that for the analogous question of $\{0,1\}$ -colorability of square-free graphs in [15], with the difference being that we directly use the constructions of 01-gadgets from the previous sections. Firstly, we know that checking $\{0,1\}$ -colorability of a $\{G_{\text{fbd}}\}$ -free graph is in NP because the problem of checking an arbitrary graph for $\{0,1\}$ -colorability is in NP [15]. Suppose we are given a graph G . The idea is to construct a new graph H which is $\{G_{\text{fbd}}\}$ -free such that the problem of $\omega(G)$ -colorability of G is equivalent to the problem of $\{0,1\}$ -colorability of H . Provided the construction is achievable in polynomial time, this gives a reduction from the $\{0,1\}$ -colorability problem to the $\omega(G)$ -colorability problem (for $\omega(G) \geq 3$) which is known to be NP-complete [27].

The construction goes as follows. Replace every vertex $v \in V(G)$ by a clique of size $\omega(G)$ in H and label the corresponding vertices $v_i \in V(H)$ for $i \in [\omega(G)]$. For every edge $(u,v) \in E(G)$, connect the corresponding vertices (u_i, v_i) by a 01-gadget $\Gamma_{(u_i, v_i)}$ in H . The exact form of the gadget $\Gamma_{(u_i, v_i)}$ is left unspecified at the moment, for the polynomial time reduction it is only important that it is finite (i.e., $|V(\Gamma_{(u_i, v_i)})|$ and $|E(\Gamma_{(u_i, v_i)})|$ are finite), so that $|V(H)| \leq \omega(G)(|V(\Gamma_{(u_i, v_i)})|_{\text{max}} - 1)|V(G)|$ and $|E(H)| \leq \omega(G)|V(G)| + |E(\Gamma_{(u_i, v_i)})|_{\text{max}}|E(G)|$, i.e., $|V(H)| = O(|V(G)|)$ and $|E(H)| = O(|E(G)| + |V(G)|)$.

We first verify that H is $\{G_{\text{fbd}}\}$ -free. We do this by showing that H is in fact faithfully realizable in dimension $\omega(G)$ and consequently free of the forbidden subgraphs for that dimension. For the vertices $v \in V(G)$, the actual representation of the vertices $v_i \in V(H)$ is chosen independent of the exact structure of the graph, i.e., for any G with $|V(G)| = n$, we choose a fixed faithful orthogonal representation $\{|v_i\rangle\}$ for $v \in V(G)$ and $i \in [\omega(G)]$. Indeed, to show the realizability of the rest of H , it suffices to show the realizability of the vertices v_1 for $v \in V(G)$, since the representation for the remaining vertices v_i for $i \geq 2$ can be readily obtained by a cyclic permutation $\Pi_i : |j\rangle \mapsto |j+i\rangle$ with the sum taken modulo $\omega(G)$. The structure of the graph is then incorporated by means of an appropriate choice of the gadgets $\Gamma_{(u_i, v_i)}$. The crucial idea behind the construction is that there exist finite sized gadgets (with faithful representations) for any two distinct vertices as shown in Prop. 4. So that for any edge $(u,v) \in E(G)$, we use a gadget $\Gamma_{(u_1, v_1)}$ from Prop. 4 (the same gadget is used for the other pairs (u_i, v_i)) corresponding to the required overlap $|\langle u_1 | v_1 \rangle|$. Now, since the representation is faithful, we do not have different vertices represented by the same vector. As such, the construction from Prop. 4 yields a finite sized gadget for any pair of vertices (u_i, v_i) .

The proof that checking $\{0,1\}$ -colorability of the $\{G_{\text{fbd}}\}$ -free graph H is equivalent to checking the $\omega(G)$ -colorability of G (which is NP-complete) follows along analogous lines to the proof in [15] and we present it here for completeness. Firstly, we show that H is $\{0,1\}$ -colorable if G is $\omega(G)$ -colorable. Consider the intermediate situation when we form a graph G' by replacing every vertex $v \in V(G)$ by a clique of size $\omega(G)$ and labeling the corresponding vertices $v_i \in V(G')$ for $i \in [\omega(G)]$. For every edge $(u,v) \in E(G)$, connect the corresponding vertices (u_i, v_i) by an edge in G' . The strategy is to show that if G is $\omega(G)$ -colorable, then G' admits a valid $\{0,1\}$ -assignment. Suppose G is $\omega(G)$ -colorable, and $c : V(G) \rightarrow [\omega(G)]$ is an optimal coloring. We define the $\{0,1\}$ -coloring of G' by

$$c'(v_i) = \begin{cases} 1, & \text{for } i = c(v) \\ 0, & \text{else} \end{cases}$$

The fact that this is a valid $\{0,1\}$ -coloring of G' follows the proof of Lemma 1 in [15]. We now derive the $\{0,1\}$ -coloring of H from that of G' by seeing that each of the gadgets in Prop. 4 can be $\{0,1\}$ -colored in all three cases, when the distinguished vertices u_i, v_i have the assignments: (i) $f(u_i) = 0, f(v_i) = 0$, (ii) $f(u_i) = 0, f(v_i) = 1$, and $f(u_i) = 1, f(v_i) = 0$. This is done by checking that such a valid $\{0,1\}$ -coloring exists for the Clifton gadget in Fig. 1 in each of the three cases. The $\{0,1\}$ -coloring can be extended to the entire gadget iteratively by following the procedure shown in the proof of Prop. 4. This gives a valid $\{0,1\}$ -coloring of H .

We now show that a valid $\{0,1\}$ -coloring of H also implies that G is $\omega(G)$ -colorable. Let $f : V(H) \rightarrow \{0,1\}$ be a valid $\{0,1\}$ assignment of H . For every $v \in V(G)$, by the fact that we have a valid $\{0,1\}$ -coloring, exactly one of the vertices $v_i \in V(H)$ is assigned value 1, i.e., $f(v_i) = 1$. One can then define a $\omega(G)$ -coloring $c : V(G) \rightarrow [\omega(G)]$ by $c(v) = i \leftrightarrow c(v_i) = 1$ for every $v \in V(G)$. It is clear that this is a valid coloring since if $(u, v) \in E(G)$ we have by the property of the gadget that at most one of u_i, v_i is assigned value 1, i.e., either $f(u_i) = 0$ or $f(v_i) = 0$. Thus, the $\{0,1\}$ -colorability of the $\{G_{\text{fbd}}\}$ -free graph H is equivalent to the $\omega(G)$ -colorability of G . From [27], we know that for $\omega(G) \geq 3$, this problem is NP-complete, which finishes the proof. \square

It is also interesting to examine the complexity of identifying 01-gadgets. In this case, it appears to be necessary to enumerate all $\{0,1\}$ -colorings of a given graph and to check $O(n^2)$ vertices to identify the possible distinguished vertices. Note that for a graph with n vertices there are 2^n possible $\{0,1\}$ -colorings so that it is not apparent whether even a polynomially checkable certificate exists for this problem. Peeters in [31] gave a polynomial time reduction preserving graph planarity of the problem of testing $\xi(G) \leq 3$ to the problem of testing whether the chromatic number $\chi(G)$ is less than or equal to 3, which is a well-known NP-complete problem, so that it is hard to check whether $d(G) \leq 3$ already for the case of planar graphs.

8 Randomness from 01-gadgets

In this section, we give a brief outline of how 01-gadgets may be linked to device-independent randomness certification. Namely, when two parties Alice and Bob perform locally the measurements from the Clifton gadget on their half of a maximally entangled state (in $\mathbb{C}^3 \otimes \mathbb{C}^3$), we will show that some specific outcome of their joint measurements has probability bounded from above and below (and this holds in all no-signaling theories). This can be inserted into a fully device-independent protocol as given in [10], the details are deferred to a separate paper [29]. To show

how the Clifton gadget can be used for randomness amplification we first consider a non-contextual assignment of probabilities to its vertices v satisfying

$$\sum_{v \in \text{clique}} p_v \leq 1, \quad \sum_{v \in \text{maximum clique}} p_v = 1 \quad (27)$$

This is the same requirement as Eq.(3), but we now assign not necessarily zeros and ones, but probabilities (i.e., values in $[0,1]$ rather than in $\{0,1\}$). Recall that such an assignment was also considered in our discussion of the extended Kochen-Specker theorem in Section 6. Now, since the gadgets are $\{0,1\}$ colorable, such an assignment of zeros and ones is possible, although in the $\{0,1\}$ assignment, it is not possible to assign 1's to both vertices 1 and 8. Here, we will first show, that even if we assign probabilities, we still cannot have $p_1 = p_8 = 1$, and we will provide a quantitative bound for this. Indeed, let us write Eq.(27) explicitly for the cliques in question from the Clifton gadget in Fig. 1:

$$\begin{aligned} p_1 + p_2 &\leq 1, & p_1 + p_6 &\leq 1, & p_4 + p_5 &\leq 1, \\ p_7 + p_8 &\leq 1, & p_3 + p_2 &\leq 1 \end{aligned} \quad (28)$$

for non-maximal cliques and

$$p_2 + p_3 + p_4 = 1, \quad p_5 + p_6 + p_7 = 1 \quad (29)$$

for the two maximum cliques. We sum up all the inequalities (28), and get

$$2p_1 + p_2 + p_3 + p_4 + p_5 + p_6 + p_7 + 2p_8 \leq 5. \quad (30)$$

Using (29) we then obtain

$$p_1 + p_8 \leq \frac{3}{2}. \quad (31)$$

To exploit this feature for randomness amplification, we consider a maximally entangled state shared by two parties. The parties will measure observables composed of the projectors given by the quantum representation (if the clique is not maximal, one simply adds a third orthogonal projector to obtain a complete measurement). Recall here that a set of eight projectors $P_v = |u_v\rangle\langle u_v|$ that is compatible with the Clifton graph is given in Eq.(5). Projectors of Alice will be denoted A_v and those of Bob B_v , and the probability of obtaining outcome v , while measuring observable containing v , will be denoted by $p(A_v = 1)$. We correspondingly denote by $p(A_v = 0)$ the probability that the outcome v was not obtained. Clearly $p(A_v = 1) + p(A_v = 0) = 1$. Now, we shall show using no-signaling (which will impose non-contextuality), that the probability $p(A_1 = 1, B_8 = 1)$ is bounded from above. To see this, we apply Eq.(31) to Alice's observables and get

$$p(A_1 = 1) + p(A_8 = 1) \leq \frac{3}{2} \quad (32)$$

From the correlations in the maximally entangled state, we have that

$$p(A_8 = 1) = p(B_8 = 1) \quad (33)$$

giving

$$p(A_1 = 1) + p(B_8 = 1) \leq \frac{3}{2}. \quad (34)$$

Now, from no-signaling we have

$$\begin{aligned} p(A_1 = 1) &= p(A_1 = 1, B_8 = 0) + p(A_1 = 1, B_8 = 1), \\ p(B_8 = 1) &= p(A_1 = 0, B_8 = 1) + p(A_1 = 1, B_8 = 1). \end{aligned} \quad (35)$$

Summing these and applying (34) we get

$$\begin{aligned} p(A_1 = 1, B_8 = 0) + p(A_1 = 1, B_8 = 1) + \\ p(A_1 = 0, B_8 = 1) + p(A_1 = 1, B_8 = 1) \leq \frac{3}{2} \end{aligned} \quad (36)$$

and hence

$$p(A_1 = 1, B_8 = 1) \leq \frac{3}{4}. \quad (37)$$

Thus we have obtained, that the probability of the event $(A_1, B_8) = (1, 1)$ is bounded from above. We have also the lower bound

$$p(A_1 = 1, B_8 = 1) = \frac{1}{3} |\langle u_1 | u_8 \rangle|^2 \geq \frac{1}{27}. \quad (38)$$

Thus

$$\frac{1}{27} \leq p(A_1 = 1, B_8 = 1) \leq \frac{3}{4} \quad (39)$$

Therefore, the outcome $(A_1, B_8) = (1, 1)$ has randomness, which can be used in a randomness amplification scheme employing the protocol of [10]. The lower bound is $\frac{1}{27}$ in noiseless conditions, and assuming we have exactly measured the specified projectors. In a real experiment, this value may be different, but if the noise is low enough it should be close to $\frac{1}{27}$. Also the upper bound, relies on perfect correlations, which in a real experiment may be imperfect. Thus in noisy conditions, we will have less stringent lower and upper bounds, though these are certifiable by statistics from the experiment. Note that crucially we have not used explicitly Bell inequalities, nor even the KS paradox. We have simply made use of the perfect correlations between the parties and the local 01-gadget structure of Alice and Bob's observables.

9 Conclusion and Open Questions

In this paper, we have shown that there exist interesting subgraphs of the Kochen-Specker graphs that we termed 01-gadgets that encapsulate the main contradiction necessary to prove the Kochen-Specker theorem. Furthermore, as a main technical contribution, we have shown that the fundamental structures identified here, lead to clean constructions of state-independent statistical proofs of the KS theorem, of

which the famous Yu and Oh proof is a particular case. The proofs given here provide a new perspective on these results, and serve as a useful tool to construct minimal KS sets, since efforts may be concentrated on the 01-gadget subgraphs. An extended notion of 01-gadgets also helped to provide simple constructive proofs of the extended Kochen-Specker theorem [22]. The gadgets enable a proof of the NP-completeness of checking $\{0, 1\}$ -colorability of graphs free from the forbidden subgraphs from Hilbert spaces of any dimension. Practically, the gadgets open up a highly important application of contextuality to practical device-independent randomness generation, which we study in a companion paper [29] where we provide an explicit device-independent protocol for randomness amplification based on [11, 10, 12] and Hardy paradoxes constructed using 01-gadgets.

An open question, is to find, for given overlap $|\langle v_1 | v_2 \rangle|$, the minimal 01-gadget and extended 01-gadget with the corresponding vertices v_1, v_2 playing the role of the distinguished vertices. An answer to this question would have applications for randomness generation from contextuality [29]. Another open question is whether all state-independent contextual graphs (including those going beyond KS sets such as that of Yu and Oh [25]) contain 01-gadgets as subgraphs, or even possibly as induced subgraphs. Finally, while it is known that in \mathbb{C}^3 KS sets cannot be constructed using rational vectors [5], it would be very interesting to study quantum realizations of 01-gadgets using rational vectors, to build statistical KS arguments and state-independent non-contextuality inequalities using these.

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